

Pruning for GNNs: Lower Complexity with Comparable Expressiveness

Dun Ma¹ Jianguo Chen^{2,4} Wenguo Yang² Suixiang Gao^{2,5}
Shengminjie Chen³

¹School of Advanced Interdisciplinary Sciences, University of Chinese Academy of Sciences

²School of Mathematical Sciences, University of Chinese Academy of Sciences

³Institute of Computing Technology, Chinese Academy of Sciences

⁴Academy of Mathematics and Systems Science, Chinese Academy of Sciences

⁵Zhongguancun Laboratory, Beijing, China

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- 1 Introduction
- 2 Pruned Message-Passing Framework
- 3 Experiment

- Standard Message-Passing GNN:

$$M_v^l = AGG^l(\{\{H_u^{l-1} | u \in N(v)\}\}) \quad (1)$$

$$H_v^l = COB^l(H_v^{l-1}, M_v^l) \quad (2)$$

- Multi-Aggregation GNN:

$$M_v^{l,k} = AGG_k^l(\{\{H_u^{l-1} | u \in N_{path}^k(v)\}\}), \quad (3)$$

$$\mathbf{M}_v^l = (M_v^{l,1}, M_v^{l,2}, \dots, M_v^{l,K}) \quad (4)$$

$$H_v^l = COB^l(H_v^{l-1}, \mathbf{M}_v^l) \quad (5)$$

Theorem

Two adjacency matrices are indistinguishable by the 1-WL test if and only if $e(A_G) = e(A_{G'})$ for all $e \in ML(\mathcal{L}_1)$, where $\mathcal{L}_1 = \{.,^\top, \mathbf{1}, \text{diag}\}$.

- Matrix languages $ML(\mathcal{L})$ can be formalised through composition of linear algebra operations. A linear algebra operation takes a number of matrices as input and returns another matrix (vector or scalar).
- For example, if A is a adjacency matrix, then $e(A) = \mathbf{1}^\top A \mathbf{1}$ is a scalar sentence in $ML(\mathcal{L})$ with $\mathcal{L} = \{.,^\top, \mathbf{1}\}$, computing the number of edges in G .
- Denote $(G, G') \in GI_{\mathcal{A}}^L$, if graph isomorphism algorithm \mathcal{A} decides (G, G') is isomorphic at L^{th} iteration. $GI_{\mathcal{A}}^L \subset GI_{\mathcal{B}}^L$ denotes \mathcal{A} more powerful than \mathcal{B} .

Algebra Operation of Matrix Language

conjugate transposition ($\text{op}(e) = e^*$)

$$e(\nu(X)) = A \in \mathbb{C}^{m \times n}$$

$$e(\nu(X))^* = A^* \in \mathbb{C}^{n \times m}$$

$$(A^*)_{ij} = \overline{A_{ji}}$$

one-vector ($\text{op}(e) = \mathbf{1}(e)$)

$$e(\nu(X)) = A \in \mathbb{C}^{m \times n}$$

$$\mathbf{1}(e(\nu(X))) = \mathbf{1} \in \mathbb{C}^{m \times 1}$$

$$\mathbf{1}_i = 1$$

diagonalization of a vector ($\text{op}(e) = \text{diag}(e)$)

$$e(\nu(X)) = A \in \mathbb{C}^{m \times 1}$$

$$\text{diag}(e(\nu(X))) = \text{diag}(A) \in \mathbb{C}^{m \times m}$$

$$\begin{aligned} \text{diag}(A)_{ii} &= A_i, \\ \text{diag}(A)_{ij} &= 0, i \neq j \end{aligned}$$

matrix multiplication ($\text{op}(e_1, e_2) = e_1 \cdot e_2$)

$$e_1(\nu(X)) = A \in \mathbb{C}^{m \times n}$$

$$e_1(\nu(X)) \cdot e_2(\nu(X)) = C \in \mathbb{C}^{m \times o}$$

$$C_{ij} = \sum_{k=1}^n A_{ik} \times B_{kj}$$

$$e_2(\nu(X)) = B \in \mathbb{C}^{n \times o}$$

scalar multiplication ($\text{op}(e) = c \times e, c \in \mathbb{C}$)

$$e(\nu(X)) = A \in \mathbb{C}^{m \times n}$$

$$c \times e(\nu(X)) = B \in \mathbb{C}^{m \times n}$$

$$B_{ij} = c \times A_{ij}$$

trace ($\text{op}(e) = \text{tr}(e)$)

$$e(\nu(X)) = A \in \mathbb{C}^{m \times m}$$

$$\text{tr}(e(\nu(X))) = c \in \mathbb{C}$$

$$c = \sum_{i=1}^m A_{ii}$$

pointwise matrix multiplication (Schur-Hadamard) ($\text{op}(e_1, e_2) = e_1 \odot e_2$)

$$e_1(\nu(X)) = A \in \mathbb{C}^{m \times n}$$

$$e_1(\nu(X)) \odot e_2(\nu(X)) = C \in \mathbb{C}^{m \times n}$$

$$C_{ij} = A_{ij} \times B_{ij}$$

$$e_2(\nu(X)) = B \in \mathbb{C}^{m \times n}$$

a_k -walk Message-Passing Framework

Theorem

Given a positive integer sequence a_k and a pair of graphs (G, G') , $S_k = \sum_{t \in [k]} a_t$, if a_k is viewable, then $\forall l \in \mathbb{N}^+ \quad Gl_{a_k\text{-walk}}^l \subseteq Gl_{WL}^{S_l}$.

- Given a positive integers sequence a_k , $S_k = \sum_{t \in [k]} a_t$. a_k is viewable if $\forall k \in \mathbb{N}^+, r \in [S_k]$ (The subset sums of a_k are dense in \mathbb{N}^+)
- a_k -walk Message-Passing GNN:

$$M_v^k = \underbrace{AGG(\cdots AGG}_{a_k \text{ times}}(\{H_u | u \in N(v)\})). \quad (6)$$

$$H_v^l = COB^l(H_v^{l-1}, M_v^l) \quad (7)$$

Pruned Multi-Aggregation GNN

Theorem

Given a pair of graphs (G, G') and $K \in \mathbb{N}^+$, $\forall L \in \mathbb{N}^+$, for K -Path GNN, the expressiveness of pruned K -Path framework is as powerful as K -Path framework: $GI_{PR\ K-P}^L = GI_{K-P}^L$. For K -Hop GNN, Pruned framework have same expressiveness referring to regular graphs and strong regular graphs: $(RG \cap GI_{PR\ 2-H}^L) \subseteq (RG \cap GI_{2-H}^L)$, $(SRG \cap GI_{RE\ K-H}^L) \subseteq (SRG \cap GI_{K-H}^L)$.

- Pruned Multi-Aggregation GNN: When $l \leq K$:

$$M_v^{l,k} = AGG_k^l(\{\{H_u^{l-1} | u \in N_{path}^k(v)\}\})(k \leq l \leq K) \quad (8)$$

$$\mathbf{M}_v^l = (M_v^{l,l}, M_v^{l,l+1}, \dots, M_v^{l,K}) \quad (9)$$

$$H_v^l = COB^l(H_v^{l-1}, \mathbf{M}_v^l) \quad (10)$$

else when $l > K$:

$$\begin{aligned} M_v^{l,K} &= AGG_K^l(\{\{H_u^{l-1} | u \in N_{path}^K(v)\}\}) \\ H_v^l &= COB^l(H_v^{l-1}, M_v^{l,K}) \end{aligned} \quad (11)$$

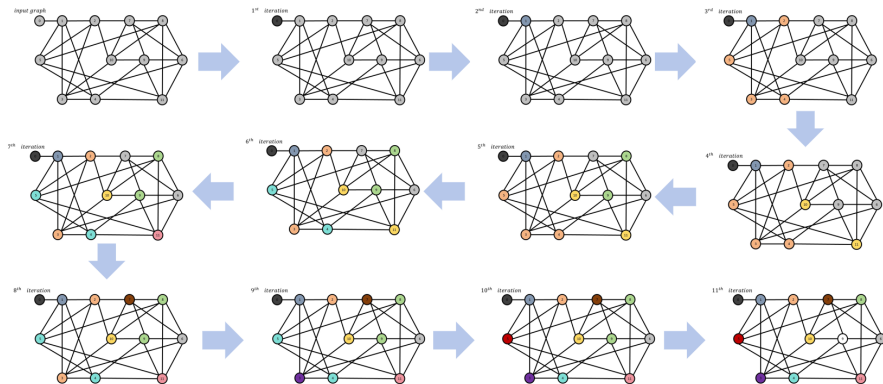
The Pruning Weisfeiler-Lehman Algorithm

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- 1: **Input:** Graph $G = (V, E)$, number of iterations L .
 - 2: **Initialization:** $\forall v \in V, \chi^0(v), l = 0$.
 - 3: **while** $l \leq L$ **do**
 - 4: $l = l + 1, t = 1$.
 - 5: $m_1^l(v) = \text{Hash}(\{\{\chi^{l-1}(u) : u \in N(v)\}\})$.
 - 6: **for** $\forall v \in V$ **do**
 - 7: **while** $t < 2^{l-1}$ **do**
 - 8: $t = t + 1$.
 - 9: $m_t^l(v) = \text{Hash}(\{\{m_{t-1}^l(u) : u \in N(v)\}\})$.
 - 10: **end while**
 - 11: $\chi^l(v) = \text{Hash}(\chi^{l-1}(v), m_{2^{l-1}}^l(v))$.
 - 12: **end for**
 - 13: **end while**
 - 14: **Output:** Final labels $\chi^L(v)$ for all $v \in V$.
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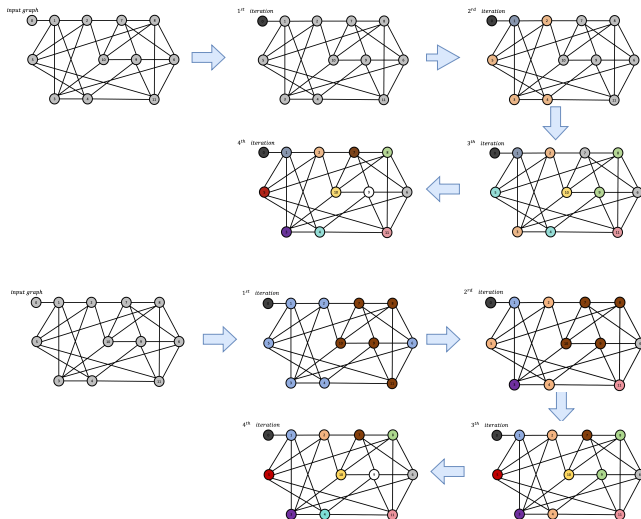
The Pruned Multi-Weisfeiler-Lehman Algorithm

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1: Input: Graph  $G = (V, E, X)$ , number of iterations  $L$ .
2: Initialization:  $\forall v \in V, \chi^0(v), l = 0$ .
3: while  $l \leq K$  do
4:    $l = l + 1$ .
5:   for  $v \in V$  do
6:     for  $t \in [l, K]$  do
7:        $\chi_t^l(v) = \text{Hash}(\{\{\chi^{l-1}(u) : u \in N_t^t(v)\}\})$ .
8:     end for
9:      $\mathbf{X}^l(v) = (\chi_l^l(v), \chi_{l+1}^l(v), \dots, \chi_K^l(v))$ .
10:     $\chi^l(v) = \text{Hash}(\chi^{l-1}(v), \mathbf{X}^l(v))$ .
11:   end for
12: end while
13: while  $l \leq L$  do
14:    $l = l + 1$ .
15:   for  $v \in V$  do
16:      $\chi^l(v) = \text{Hash}(\chi^{l-1}(v), \{\{\chi^{l-1}(u) : u \in N_{path}^L(v)\}\})$ .
17:   end for
18: end while
19: Output: Final labels  $\chi^L(v)$  for all  $v \in V$ .
```

Process of WL Test on Long-refinement Graph



Process of Pruned WL Test and Pruned 2-hop WL Test on Long-refinement Graph



Comparison on training efficiency

Table: Comparison on training efficiency

Model	COLLAB		NCI1		IMDB-B		IMDB-M		MUTAG		PROTEINS	
	Time	Acc (%)	Time	Acc (%)	Time	Acc (%)	Time	Acc (%)	Time	Acc (%)	Time	Acc (%)
GIN(3)	1.104	74.8 \pm 1.3	0.480	71.9 \pm 0.5	0.251	71.9 \pm 0.3	0.304	49.9 \pm 0.0	0.889	89.4 \pm 0.4	0.268	73.7 \pm 0.7
PR GIN(1)	1.060	73.9 \pm 0.0	0.461	72.9 \pm 1.4	0.209	69.9 \pm 2.0	0.284	50.6 \pm 0.3	0.886	88.5 \pm 0.0	0.233	72.2 \pm 1.9
GIN(7)	1.638	77.4 \pm 1.6	0.748	71.5 \pm 1.4	0.578	72.6 \pm 0.3	0.534	51.1 \pm 0.3	0.904	89.4 \pm 1.0	0.527	76.3 \pm 0.2
PR GIN(124)	1.284	76.4 \pm 0.7	0.6961	75.4 \pm 0.2	0.481	71.7 \pm 1.4	0.464	52.0 \pm 0.5	0.916	92.0 \pm 0.4	0.425	74.1 \pm 1.0
GIN(10)	2.142	74.7 \pm 0.6	1.122	75.9 \pm 1.3	0.948	72.1 \pm 2.8	0.898	49.7 \pm 1.5	0.874	87.7 \pm 0.2	0.867	72.3 \pm 0.0
PR GIN(1234)	1.689	75.6 \pm 0.3	0.981	74.7 \pm 0.2	0.710	71.5 \pm 0.5	0.780	51.2 \pm 0.9	0.929	90.7 \pm 2.1	0.667	72.5 \pm 2.2
2-Hop(3)	1.357	76.8 \pm 0.8	0.608	73.6 \pm 0.9	0.410	71.0 \pm 0.7	0.415	50.1 \pm 1.5	0.910	91.0 \pm 0.0	0.394	69.5 \pm 1.3
PR 2-Hop(3)	1.180	75.1 \pm 1.1	0.528	76.5 \pm 1.6	0.357	71.5 \pm 0.5	0.361	52.5 \pm 0.7	0.929	91.3 \pm 1.5	0.342	73.3 \pm 0.3
2-Hop(5)	1.927	74.2 \pm 0.5	0.880	70.6 \pm 1.7	0.680	68.8 \pm 0.8	0.628	49.5 \pm 0.6	0.894	88.3 \pm 1.0	0.621	71.0 \pm 0.8
PR 2-Hop(5)	1.606	74.5 \pm 0.4	0.734	71.1 \pm 1.5	0.566	69.7 \pm 0.2	0.523	48.0 \pm 2.2	0.871	88.7 \pm 1.5	0.517	72.0 \pm 0.3
2-Path(3)	1.385	75.6 \pm 0.0	0.620	73.0 \pm 0.8	0.418	71.4 \pm 0.1	0.423	49.5 \pm 1.8	0.9045	88.7 \pm 1.7	0.402	72.7 \pm 2.0
PR 2-Path(3)	1.385	76.1 \pm 0.3	0.632	75.5 \pm 0.4	0.488	74.2 \pm 1.9	0.451	51.6 \pm 0.4	0.915	91.6 \pm 0.0	0.446	76.9 \pm 0.7
2-Path(5)	2.111	76.0 \pm 0.2	0.964	74.2 \pm 0.3	0.745	70.0 \pm 0.6	0.688	51.3 \pm 0.1	0.890	89.4 \pm 0.4	0.680	69.3 \pm 1.0
PR 2-Path(5)	1.730	72.1 \pm 2.0	0.790	70.5 \pm 2.3	0.610	71.8 \pm 1.5	0.564	50.2 \pm 0.2	0.898	88.9 \pm 0.8	0.557	71.9 \pm 2.4

Thank You for Listening!