

Statistical and Computational Guarantees of Kernel Max-Sliced Wasserstein Distances

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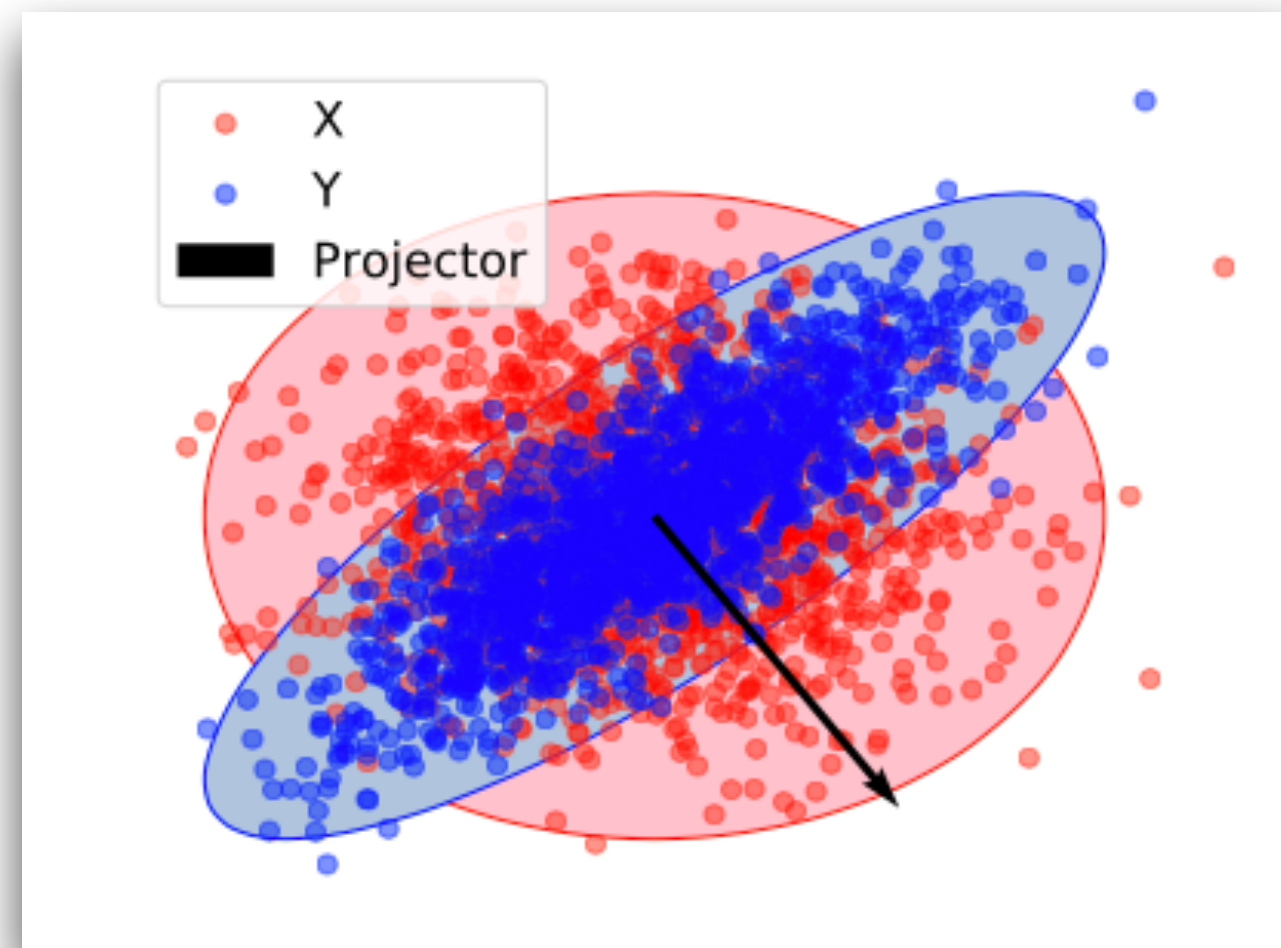
Background of Hypothesis Testing

❖ **Given:** *high-dimensional* samples from unknown distributions μ and ν

❖ **Goal:** determine whether μ and ν differ

- Project data linearly and follow by Wasserstein testing [1]

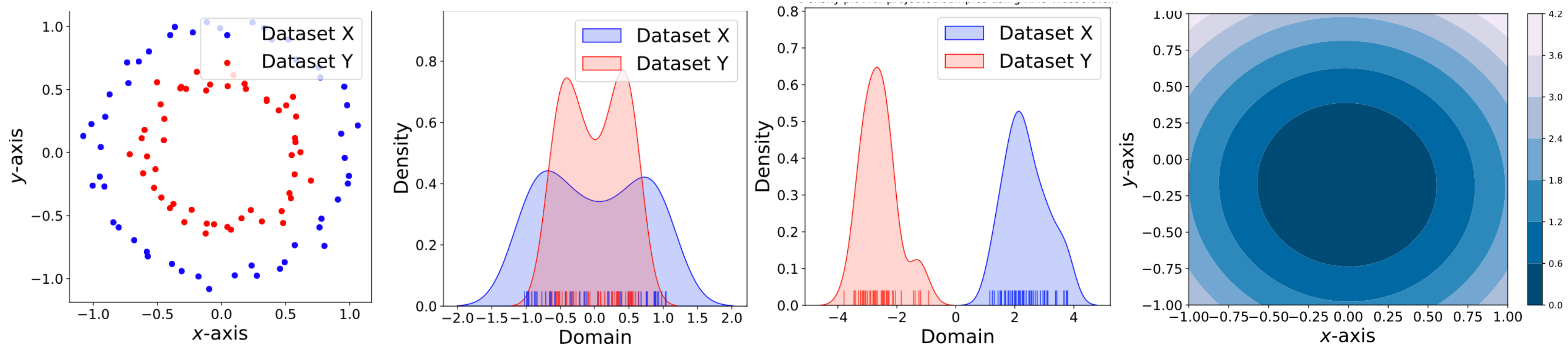
$$\text{MS}(\mu, \nu) = \max_{f: f(x)=a^\top x} \mathbf{W}(f_\# \mu, f_\# \nu)$$



Testing with Nonlinear Dimensionality Reduction

- Use **nonlinear** operator modeled by kernel method

$$\mathbf{KMS}(\mu, \nu) = \max_{f \in \mathcal{H}: \|f\|_{\mathcal{H}} \leq 1} \mathbf{W}(f_{\#}\mu, f_{\#}\nu)$$



(a) Scatter plot of data

(b) Density of projected samples using MS

(c) Density of projected samples using KMS

(d) Heatmap of estimated nonlinear projector

Kernel Max-Sliced Wasserstein Distance

$$\mathbf{KMS}(\mu, \nu) = \max_{f \in \mathcal{H}: \|f\|_{\mathcal{H}} \leq 1} \mathbf{W}(f_{\#}\mu, f_{\#}\nu)$$

- \mathcal{H} : reproducing kernel Hilbert space (RKHS) dependent on $k(\cdot, \cdot) : \mathbb{R}^{D \times D} \rightarrow \mathbb{R}^1$
- $f : \mathbb{R}^D \rightarrow \mathbb{R}^1$: nonlinear projector

1. Statistical guarantees of KMS
2. Computational guarantees of KMS
3. Practical Applications

Finite-Sample Guarantees

$$\mathbf{KMS}(\mu, \nu) = \max_{f \in \mathcal{H}: \|f\|_{\mathcal{H}} \leq 1} W(f_{\#}\mu, f_{\#}\nu)$$

Theorem (Informal). Assume $k(x, x) \leq A, \forall x$. With high probability,

$$\mathbf{KMS}(\mu, \hat{\mu}_n) = \mathcal{O}(n^{-1/2}).$$

- $\mathcal{O}(\cdot)$ hides constant depending on A
- **KMS** breaks the curse of dimensionality of Wasserstein distance
- Free of distribution assumptions
 - MS distance [Sloan N et. al, 2022, Tianyi et. al, 2021]: finite diameter of support
 - KMS distance [Wang et. al, 2022]: light-tailed distribution

Finite-Dimensional Reformulation

$$\text{KMS}(\hat{\mu}_n, \hat{\nu}_n) = \max_{\omega \in \mathbb{R}^{2n}: \|\omega\|_2=1} \left\{ \min_{\pi \in \Gamma_n} \sum_{i,j=1}^n \pi_{i,j} (M_{i,j}^\top \omega)^2 \right\}$$

- $M_{i,j} \in \mathbb{R}^{2n}$: concatenation of kernel valued on data points
- Non-concave quadratic optimization problem

Theorem. $\text{KMS}(\hat{\mu}_n, \hat{\nu}_n)$ is \mathcal{NP} -hard to compute

(KMS Problem)

\supseteq

$$\max_{\omega: \|\omega\|_2=1} \min_{i \in [n]} \omega^\top A_i A_i^\top \omega$$

(Fair-PCA with rank-1 data)

\supseteq

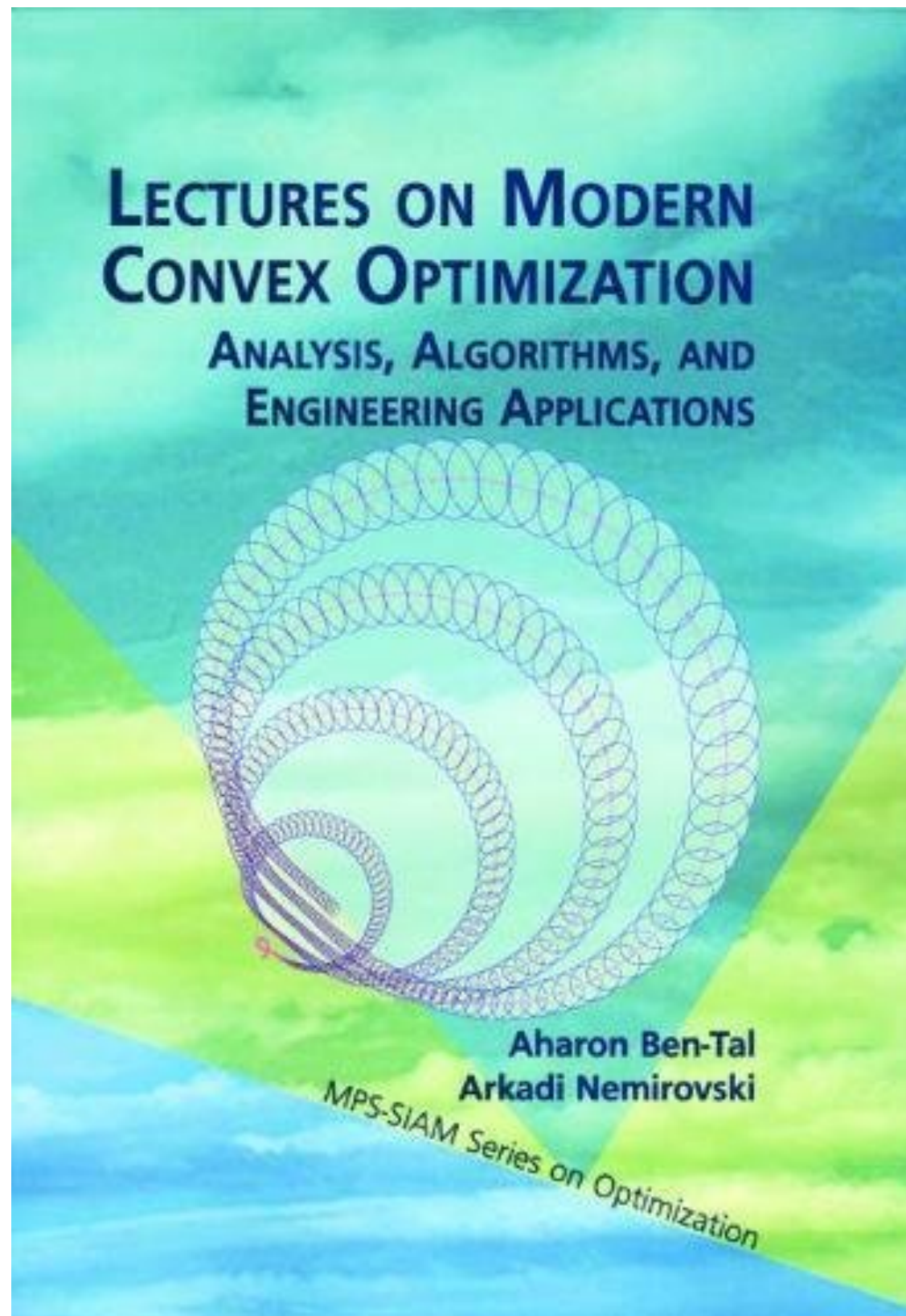
Give integers a_1, \dots, a_N ,
determine whether binary
variables $\{x_i\}_{i=1}^N \in \{-1, 1\}^N$
exist such that $\sum_{i=1}^N a_i x_i = 0$?

(Partition)

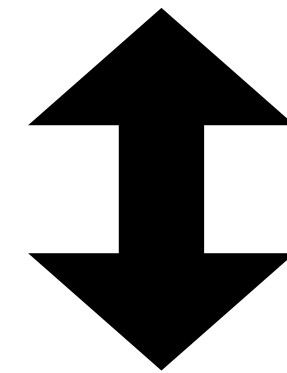
Finite-Dimensional Reformulation

$$\text{KMS}(\hat{\mu}_n, \hat{\nu}_n) = \max_{\omega \in \mathbb{R}^{2n}: \|\omega\|_2=1} \left\{ \min_{\pi \in \Gamma_n} \sum_{i,j=1}^n \pi_{i,j} \langle M_{i,j} M_{i,j}^\top, \omega \omega^\top \rangle \right\}$$

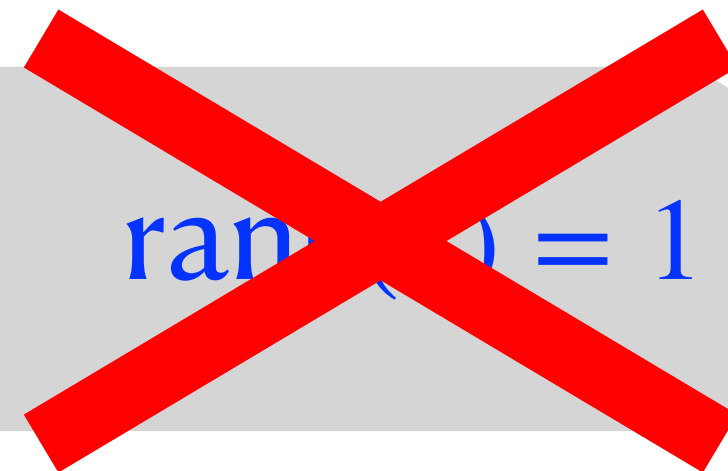
- Approximation algorithm using semidefinite relaxation (SDR):



$$S = \omega \omega^\top, \quad \omega \in \mathbb{R}^{2n}, \quad \|\omega\|_2 = 1$$



$$S \succeq 0, \quad \text{Trace}(S) = 1, \quad \text{rank}(S) = 1$$



Semidefinite Relaxation (SDR)

$$\begin{aligned} \text{KMS}(P_n, Q_n) &= \max_{S \geq 0, \text{Trace}(S)=1, \text{rank}(S)=1} \left\{ \min_{\pi \in \Gamma_n} \sum_{i,j=1}^n \pi_{i,j} \langle M_{i,j} M_{i,j}^\top, S \rangle \right\} \\ &\leq \max_{S \geq 0, \text{Trace}(S)=1} \left\{ \min_{\pi \in \Gamma_n} \sum_{i,j=1}^n \pi_{i,j} \langle M_{i,j} M_{i,j}^\top, S \rangle \right\} \end{aligned}$$

Theorem (Informal). Stochastic gradient method with biased oracles solves SDR up to δ optimality gap with operational complexity

$$\tilde{\mathcal{O}}(n^3 \delta^{-3})$$

Quality of Semidefinite Relaxation

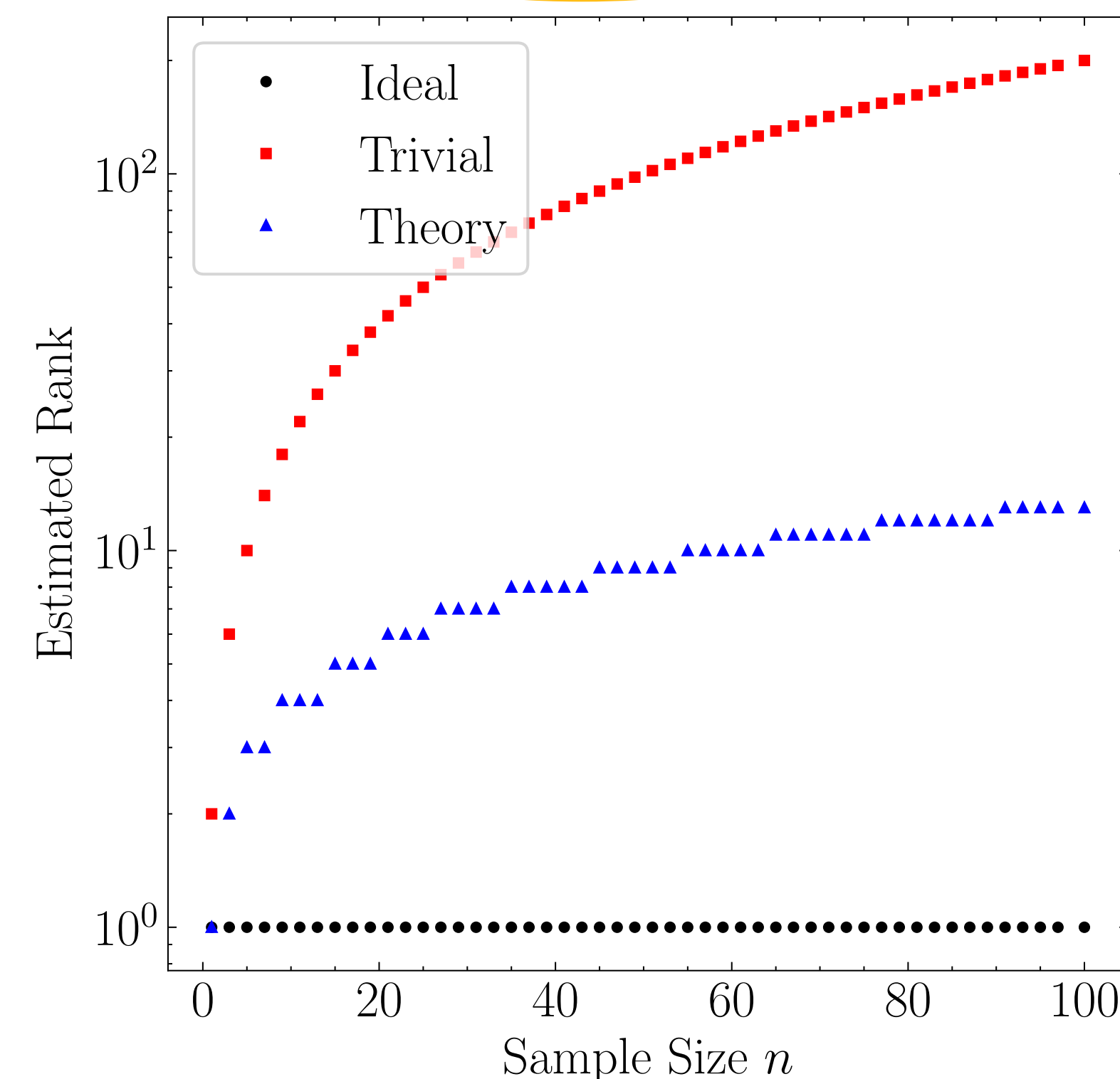
- (KMS) = $\max_{S \geq 0, \text{Trace}(S)=1, \text{rank}(S)=1} \left\{ \min_{\pi \in \Gamma_n} \sum_{i,j=1}^n \pi_{i,j} \langle M_{i,j} M_{i,j}^\top, S \rangle \right\}$
- (SDR) = $\max_{S \geq 0, \text{Trace}(S)=1} \left\{ \min_{\pi \in \Gamma_n} \sum_{i,j=1}^n \pi_{i,j} \langle M_{i,j} M_{i,j}^\top, S \rangle \right\}$

Smaller rank of optimal solution yields better performance

Theorem. There exists an optimal solution to (SDR) with

$$\text{rank } k \triangleq 1 + \left\lfloor \sqrt{2n + \frac{9}{4}} - \frac{3}{2} \right\rfloor.$$

- $(\text{KMS}) \leq (\text{SDR}) \leq \max_{S \geq 0, \text{Trace}(S)=1, \text{rank}(S)=k} \left\{ \min_{\pi \in \Gamma_n} \sum_{i,j=1}^n \pi_{i,j} \langle M_{i,j} M_{i,j}^\top, S \rangle \right\}$



Summary

- A novel **non-parametric** metric for comparing high-dimensional distributions
- **Sharp** finite-sample guarantees
- Computational Guarantees:
 - A. Non-concave quadratic maximization problem: *\mathcal{NP}* -**hard**
 - B. Approximation algorithm **with performance guarantees**:
- Practical Applications:
 1. High-dimensional Two-Sample Testing
 2. Change-Point Detection