

Improved Online Confidence Bounds for Multinomial Logistic Bandits

(ICML 2025)

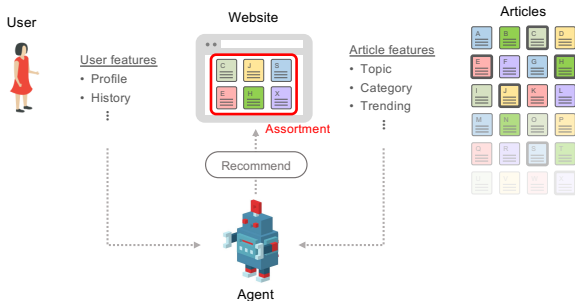
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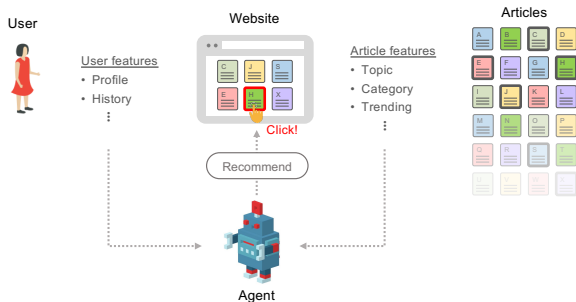
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Sequential Assortment Selection Problem



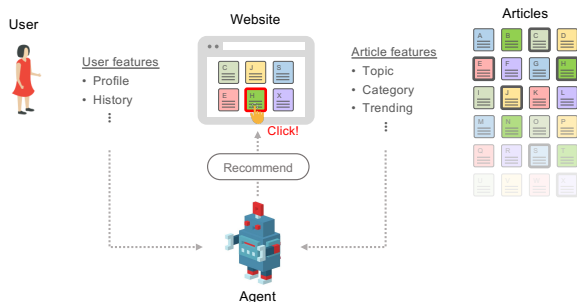
- Agent recommends an **assortment** (a set of items)

Sequential Assortment Selection Problem



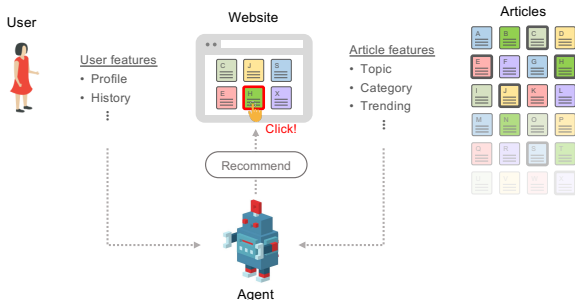
- Agent recommends an assortment (a set of items)
- User **chooses one item** from offered multiple options

Sequential Assortment Selection Problem



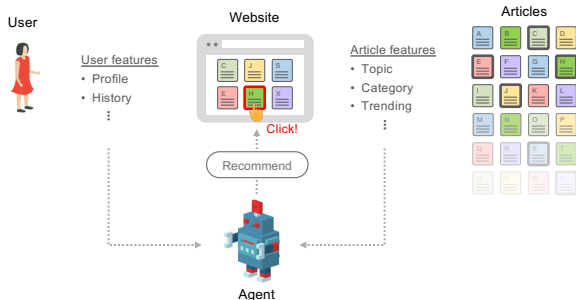
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- User chooses one item from offered multiple options
- For every round $t = 1, \dots, T$:
 1. Observe contexts $x_{ti} \in \mathbb{R}^d$ and rewards $r_{ti} \in [0, 1]$ for every item $i \in [N]$

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 3. Observe the user click decision $c_t \in S_t \cup \{0\}$ ("0": outside option)

Multinomial Logit (MNL) Choice Model (McFadden, 1977)

- Probability of choosing any item i in assortment S_t :

$$p_t(i|S_t, \mathbf{w}^\star) := \frac{\exp(x_{ti}^\top \mathbf{w}^\star)}{1 + \sum_{j \in S_t} \exp(x_{tj}^\top \mathbf{w}^\star)}$$

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- Goal: Minimize $\mathbf{Reg}_T(\mathbf{w}^\star) = \sum_{t=1}^T R_t(S_t^\star, \mathbf{w}^\star) - R_t(S_t, \mathbf{w}^\star)$

Previous Works & Main Contributions

Table. T : total rounds, d : feature dimension, K : maximum assortment size, B : bound on the parameter norm, $\mathbf{1}/\kappa = \mathcal{O}(K^2 e^B)$, $\kappa_t^* = \sum_{i \in S_t^*} p_t(i|S_t^*, \mathbf{w}^*) p_t(0|S_t^*, \mathbf{w}^*) \leq 1$,

$\sigma_t^2 \leq 1$: reward variance at round t .

	Regret	Rewards	Comput. per Round
Chen et al. (2020) (MLE-UCB)	$\mathcal{O}(Bd \log(KT)\sqrt{T})$	$r_{ti} \in [0, 1]$	Intractable
Oh and Iyengar (2021) (UCB-MNL)	$\mathcal{O}(\frac{1}{\kappa} d \log T \sqrt{T}) = \mathcal{O}(K^2 e^B d \log T \sqrt{T})$	$r_{ti} \in [0, 1]$	$\mathcal{O}(t)$
Perivier and Goyal (2022) (OFU-MNL)	$\mathcal{O}(BKd \log(KT) \sqrt{\sum_{t=1}^T \kappa_t^*})$	$r_{ti} \equiv 1$	Intractable
Lee and Oh (2024) (OFU-MNL+)	$\mathcal{O}(B^{3/2} d \log K (\log T)^{3/2} \sqrt{T})$	$r_{ti} \in [0, 1]$	$\mathcal{O}(1)$

1. No results free of poly(B) dependence
2. No results completely independent of K
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This work (OFU-MNL++)	$\mathcal{O}(d \log T \sqrt{\sum_{t=1}^T \sigma_t^2})$, for large t	$r_{ti} \in [0, 1]$	$\mathcal{O}(1)$

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**First poly(B), K -free, and variance-dependent regret bound
in MNL bandits!**

Parameter Update using Online Mirror Descent (OMD)

- **MNL loss function:** For the choice response $y_{ti} = \mathbb{1}(c_t = i)$, we define:

$$\ell_t(\mathbf{w}) := - \sum_{i \in S_t} y_{ti} \log p_t(i|S_t, \mathbf{w}).$$

- **Online parameter estimation:** Estimate $\mathbf{w}^*$ by an online mirror descent (Orabona, 2019).

$$\mathbf{w}_{t+1} = \arg \min_{\mathbf{w} \in \mathcal{W}} \langle \nabla \ell_t(\mathbf{w}_t), \mathbf{w} \rangle + \frac{1}{2\eta} \|\mathbf{w} - \mathbf{w}_t\|_{\tilde{H}_t}^2, \quad \forall t \geq 1.$$

- ▶ Step-size parameter: $\eta > 0$
- ▶ Hessian matrices: $H_t := \lambda \mathbf{I}_d + \sum_{s=1}^{t-1} \nabla^2 \ell_s(\mathbf{w}_{s+1})$ and $\tilde{H}_t := H_t + \eta \ell_t(\mathbf{w}_t)$

- Computation cost: $\mathcal{O}(Kd^3)$
- Storage cost: $\mathcal{O}(d^2)$

Result 1: Improved Online Confidence Bound

Theorem 1. Improved online confidence bound

Let $\lambda > 0$ and $\mathcal{T} \subseteq [T]$ denote the set of update rounds. For all $t \in \mathcal{T}$, we assume the following update conditions hold:

$$\sup_{\mathbf{w} \in \mathcal{W}_t} |x_{ti}^\top (\mathbf{w} - \mathbf{w}^*)| \leq \alpha, \quad \forall i \in S_t,$$

where $\mathcal{W}_t$ is a compact convex set. Then, with high probability, we have:

$$\|\mathbf{w}_t - \mathbf{w}^*\|_{H_t} = \mathcal{O}\left(\alpha\sqrt{d\log(t/\delta)} + \alpha\sqrt{\lambda}\right).$$

- Updating at every round ($\mathcal{T} = [T]$) yields $\mathcal{O}(B\sqrt{d\log t})$.
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- **Key implication:** If $\mathcal{W}_t$ is constructed so that α stays small, and updates occur only in this case, we obtain $\mathcal{O}(\sqrt{d\log t})$ for large t .
 - ▶ No dependence on B or K !

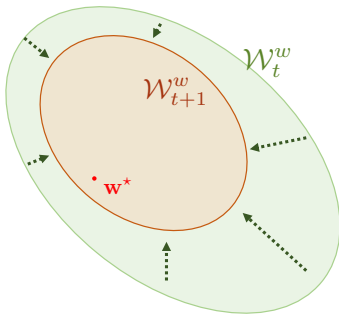
Algorithm

Algorithm OFU-MNL++

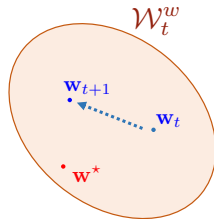
- 1: **Initialize:** $\mathcal{W}_1^w(\delta) = \mathcal{W}$, $H_1 = \lambda \mathbf{I}_d$, $H_1^w = \lambda_1^w \mathbf{I}_d$, $\mathbf{w}_1, \mathbf{w}_1^w \in \mathcal{W}$
 - 2: **for** round $t = 1, \dots, T$ **do**
 - 3: Observe feature set $\mathcal{X}_t = \{x_{ti}\}_{i=1}^N$ and rewards $\{r_{ti}\}_{i=1}^N$
 - 4: **if** $\max_{x \in \mathcal{X}_t} \|x\|_{(H_t^w)^{-1}} \geq \frac{1}{6\sqrt{2}\zeta_t(\delta)}$ **then** $\triangleright$ *Large α*
 - 5: Offer $S_t = \{i_t\}$, where $x_{ti_t} = \arg \max_{x \in \mathcal{X}_t} \|x\|_{(H_t^w)^{-1}}^2$
 - 6: Update $(\mathbf{w}_{t+1}^w, H_{t+1}^w) \leftarrow \text{OMD}(\mathcal{W}, \ell_t, H_t^w, \mathbf{w}_t^w, \eta^w)$
 - 7: Calculate $\mathcal{W}_{t+1}^w(\delta) \leftarrow \left\{ \mathbf{w} \in \mathbb{R}^d \mid \|\mathbf{w} - \mathbf{w}_{t+1}^w\|_{H_{t+1}^w} \leq \zeta_{t+1}(\delta) \right\}$
 - 8: Update $H_{t+1} \leftarrow H_t$ and $\mathbf{w}_{t+1} \leftarrow \mathbf{w}_t$.
 - 9: **else** $\triangleright$ *Small α*
 - 10: Offer $S_t = \arg \max_{S \in \mathcal{S}} \tilde{R}_t(S)$
 - 11: Update $(\mathbf{w}_{t+1}, H_{t+1}) \leftarrow \text{OMD}(\mathcal{W}_t^w(\delta), \ell_t, H_t, \mathbf{w}_t, \eta)$
 - 12: Update $H_{t+1}^w \leftarrow H_t^w$, $\mathbf{w}_{t+1}^w \leftarrow \mathbf{w}_t^w$, and $\mathcal{W}_{t+1}^w(\delta) \leftarrow \mathcal{W}_t^w(\delta)$
 - 13: **end if**
 - 14: **end for**
-

Algorithm

i) Large α : update search space



ii) Small α : update search space



Result 2: Regret Bound

Theorem 2. $\text{poly}(B)$, K -free variance-dependent regret bound

With high probability, the regret of OFU-MNL++ satisfies:

$$\mathbf{Reg}_T(\mathbf{w}^\star) = \mathcal{O} \left(\left(d \log T + Bd\sqrt{\log T} \right) \sqrt{\sum_{t=1}^T \sigma_t^2} \right).$$

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- For sufficiently large T : $\mathcal{O} \left(d \log T \sqrt{\sum_{t=1}^T \sigma_t^2} \right)$.
 - ▶ First variance-dependent and $\text{poly}(B)$, K -free regret
 - ▶ Compared to Lee and Oh (2024): $\mathcal{O} \left(B^{3/2} d \log K (\log T)^{3/2} \sqrt{T} \right)$
 $\implies$ ours improves by a factor of $\mathcal{O} \left(B^{3/2} \log K \sqrt{\log T} \right)$.

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 - ▶ Since $\sigma_t^2 \leq 1$, this represents a strict improvement over $\sqrt{T}$.

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