TURF: A Two-factor, Universal, Robust, Fast Distribution Learning Algorithm

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ICML, 2022



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- Piecewise polynomial approximation to construct f^{est}



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- Since $||u p||_1 = 2$ any f^{est} suffers either $||f^{est} u||_1 \ge 1$ or $||f^{est} p||_1 \ge 1$

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Slides at: <u>www.vaishakhr.com/turf.pdf</u>

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- We call it a c-factor approximation for $\mathscr{P}_{t.d}$

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- Similarly Gaussian and their mixtures

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- Previous works achieve c' = 3c. We achieve the optimal c' = c

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 $\mathscr{P}_{t,d}$: t - # pieces, d - degree



Thank you!