# Consensus Multiplicative Weights Update: Learning to Learn using Projector-based Game Signatures

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- 1. Consensus Multiplicative Weights Update: new algorithm with local last-iterate convergence guarantees to Nash Equilibria for zero-sum bimatrix games.
  - generalizes consensus optimization from the unconstrained to the constrained case.
- 2. Learning the coefficients of the above update rule for general bimatrix games (non zero-sum).
  - can we learn last-iterate convergence in 2 Player games using a featurization of such games, the game signature?

- constrained case: *x*, *y* players' mixed strategies in probability simplex.
- 2P bimatrix game, payoffs  $f_1(x, y) = x^T A y$ ,  $f_2(x, y) = x^T B y$ .
  - First focus on the zero-sum case B = -A
- Existing algorithms displaying last-iterate convergence in the zero-sum case, with constant step-size:
  - Optimistic Multiplicative Weights Update,
  - Optimistic mirror descent with various projection methods.

- Naively adding to the gradient the 2<sup>nd</sup> order term AA<sup>T</sup> as in the unconstrained case does **not** work!
  - Nash Equilibria are not fixed points
  - A is typically not invertible (Rock-Paper-Scissors)
    - problematic regarding the Jacobian spectral radius condition

## **Consensus Multiplicative Weights Update**

- Define simplex-Hessians  $H_x := A^T \operatorname{diag}(x)A$ ,  $H_y := \operatorname{Adiag}(y)A^T$ .
  - spectrum "rescaling"  $sp(H_x) = sp(diag(x)AA^T)$ .
- simplex-Hessians make Nash Equilibria be fixed points
- Consensus Multiplicative Weights Update (CMWU) update:
  - h learning rate,  $\epsilon$  Hessian coefficient.
  - Multiplicative Weights projection P

$$x_{t+1} = P_{x_t} (hAy_t - h\epsilon H_{y_t} x_t)$$
  
$$y_{t+1} = P_{y_t} (-hA^T x_t - h\epsilon H_{x_t} y_t)$$

- What about invertibility of  $H_x$  and  $H_y$ ?
  - rock-paper-scissors still not invertible

# **Consensus Multiplicative Weights Update**

• Definition: a matrix A is weakly V-invertible if:

 $X \in V \Rightarrow AX \neq 0$ , i.e.  $Ker(A) \cap V = \{0\}$ 

- · coincides with classical invertibility when A square and  $V = \mathbb{R}^d$
- We require A to be weakly Z-invertible
  - Z space of vectors that sum to 0
  - Rock-Paper-Scissors is weakly Z-invertible.
- CMWU has local convergence guarantees in the zero-sum case and performs competitively compared to existing methods.
  - new proof technique based on eigenvalue perturbation

# Learning Last-iterate convergence in Games

- Can we learn the update rule coefficients h and  $\epsilon$ ?
- Recent research shows that the update rule needs to depend on the *nature* of the game (Cheung & Piliouras, 2020).
  - we need to encode/featurize the game
- We allow different coefficients for the 2 players
  - 4 coefficients in total

• RL problem over trajectory  $\tau$ , minimize the distance to Nash  $\delta$ :

$$\min_{\pi} \mathbb{E}_{\pi} \sum_{t=1}^{\tau} \frac{\delta(x^t, y^t)}{\delta(x^0, y^0)}$$

- MDP state: game signature + trajectory information of both players: distances to Nash, gradient, Hessian, payoffs.
- Policy actions.  $a_t = (h_1, h_2, h_1\epsilon_1, h_2\epsilon_2)$ .
- Train on mixtures of 3 components

# Game decomposition

- How to featurize/encode a 2P game?
- Use linear operators  $\rho$  that satisfy  $\rho^2 = \rho$ , a.k.a. projectors.
- The whole space is the direct sum of the Kernel and Range of any projector:

2P Games = 
$$Ker(\rho) \oplus Range(\rho)$$

• For example,  $\rho_Z$  computes the zero-sum component of any 2P game, and Id  $-\rho_Z$  its cooperative component

$$\mathbf{\rho}_{\mathbf{Z}}: (f_1, f_2) \to \frac{1}{2}(f_1 - f_2, f_2 - f_1)$$

# Game decomposition

• Given *n* commutative projectors we have:

$$2P \text{ Games} = \bigoplus_{C_i \in \{K_{\rho_i}, R_{\rho_i}\}} \bigcap_{i=1}^n C_i \quad (2^n \text{ components})$$

- This view generalizes cyclic/transitive games in (Balduzzi, 2019) from the zero-sum symmetric case to any 2P game, and unifies it with trivial games of (Cheung & Piliouras, 2020).
- Application to the case *n* = 3 gives 8 components associated to zero-sum/cooperative, symmetric/antisymmetric, transitive/cyclic games.

## Learning to converge to Nash Equilibria

The RL policy is able to exploit the game signature across a wide range of game types.

Game Type	Avg. Distance to Nash
ZT	6
ZCy	4
СТ	4
ССу	4
Avg. Mixtures of 1	5
ZT + ZCy	7
ZT + CT	3
ZT + CCy	6
ZCy + CT	15
ZCy + CCy	10
CT + CCy	5
Avg. Mixtures of 2	8
ZT + ZCy + CT	12
ZT + ZCy + CCy	19
ZT + CT + CCy	5
ZCy + CT + CCy	20
ZT + ZCy + CT + CCy	15
Avg. Mixtures of 3-4	14
Avg. k-Mixtures	9

## Learnt coefficients - Pure components



**Figure 1:** Learnt coefficients of gradient  $G_k = h_k$  and Hessian  $H_k = -h_k \epsilon_k$  of players *k* across 8 pure game components. Mean value (left). Pairwise correlations (right).

- Gradient stronger in transitive games, Hessian stronger in cyclic games.
- mirror behavior between zero-sum and cooperative games.

### Coefficient trajectory Shape - Pure Game components



**Figure 2:** Average standardized per-episode-trajectory ("shape") of coefficients of gradient  $G_k = h_k$  and Hessian  $H_k = -h_k \epsilon_k$  of players k = 1, 2 across game types, as a function of time *t*.

#### SHAP analysis



**Figure 3:** SHAP importance with respect to **(top)** gradient learning rates  $G_k = h_k$  and **(bottom)** Hessian learning rates  $H_k = -h_k \epsilon_k$  for players k for the four input groups and eight unique game types.