Analysis of stochastic Lanczos quadrature for spectrum approximation
Tyler Chen, Thomas Trogdon, and Shashanka Ubaru


Tyler Chen


Thomas Trogdon


Shashanka Ubaru






Goal: approximate spectrum

Traditionally, approximate a few eigenvalues to high accuracy.

Local spectrum approximation via power method, vanilla Lanczos, etc.

CESM contains all information about the spectrum

CESM contains all information about the spectrum and gives a natural target for a global approximation

Goal: approximate spectrum globally

The Wasserstein distance between distribution functions $\mu$ and $\nu$ is

$$
d_{\mathrm{W}}(\mu, \nu)=\int|\mu(x)-\nu(x)| \mathrm{d} x
$$

Goal: Find approximation to $\Phi$ with small Wasserstein distance, accessing A only through matrix-vector products.

The cumulative emperical spectral measure (CESM) of $\mathbf{A}$ is the distribution function

$$
\Phi(x)=\sum_{j=1}^{n} \frac{1}{n} \mathbb{1}\left[\lambda_{j} \leq x\right] .
$$

The weighted CESM of $\mathbf{A}$ and $\mathbf{v}_{i}$ is the distribution function

$$
\Psi_{i}(x)=\sum_{j=1}^{n}\left[w_{i}\right]_{j} \mathbb{1}\left[\lambda_{j} \leq x\right]
$$

where $\left[w_{i}\right]_{i}$ is the square of the projection of $\mathbf{v}_{i}$ onto the $i$-th eigenvector of $\mathbf{A}$.

The SLQ output is the average of the Gaussian quadrature approximations to the weighted CESMs,

$$
\left\langle\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}\right\rangle:=\sum_{i=1}^{n_{\mathrm{v}}} \frac{1}{n_{\mathrm{v}}}\left[\Psi_{i}\right]_{k}^{\mathrm{gq}} .
$$

Questions: For good Wasserstein approximation

- how many samples do we need?
- how many Lanczos iterations for each sample?


$$
\Phi(x)=\sum_{j=1}^{n} \frac{1}{n}\left[\mid \lambda_{j} \leq x\right]
$$



$$
\Phi(x)=\sum_{j=1}^{n} \frac{1}{n} \mathbb{1}\left[\lambda_{j} \leq x\right] \quad \Psi_{i}(x)=\sum_{j=1}^{n}\left[w_{i}\right]_{j} \mathbb{1}\left[\lambda_{j} \leq x\right]
$$



$$
\Phi(x)=\sum_{j=1}^{n} \frac{1}{n} \mathbb{1}\left[\lambda_{j} \leq x\right] \quad \Psi_{i}(x)=\sum_{j=1}^{n}\left[w_{i}\right]_{j} \mathbb{1}\left[\lambda_{j} \leq x\right] \quad\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}(x)=\sum_{j=1}^{k}\left[d_{i}\right]_{j} \mathbb{1}\left[\left[\theta_{i}\right]_{j} \leq x\right]
$$

The main cost of SLQ is running Lanczos on $\left(\mathbf{A}, \mathbf{v}_{i}\right)$

If we run $k$ iterations of Lanczos on $n_{v}$ vectors, total number of matrix vector products is $k n_{v}$.

Lanczos requires $O(n)$ storage without reorthogonalization and $O(k n)$ with.

Our analysis assumes an implementation of Lanczos which is close to exact arithmetic; for instance using reorthogonalization.

There is evidence SLQ still seems to work without reorthogonalization, but our analysis isn't necessarily directly applicable in these cases.

Lanczos (and therefore SLQ) only requires access to A through matrix-vector products!
$\mathbb{1}[$ true $]=1 \quad \mathbb{1}[$ false $]=0$

$$
\Phi(x)=\sum_{j=1}^{n} \frac{1}{n} \mathbb{1}\left[\lambda_{j} \leq x\right]=n^{-1} \operatorname{tr}(\mathbb{1}[\mathbf{A} \leq x])
$$

$$
\Psi_{i}(x)=\sum_{j=1}^{n}\left[w_{i}\right]_{j} \mathbb{1}\left[\lambda_{j} \leq x\right]=\mathbf{v}_{i}^{\top} \mathbb{1}[\mathbf{A} \leq x] \mathbf{v}_{i}
$$

## SLQ overview:

- sample $\mathbf{v}_{1}, \mathbf{v}_{2}, \ldots, \mathbf{v}_{n_{\mathrm{v}}}$ iid from unit hypersphere - for each $\mathbf{v}_{i}$, approximate weighted CESM $\Psi_{i}$ by $\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}$, obtained via $k$ iterations of Lanczos
- Output average $\left\langle\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}\right\rangle=\sum_{i=1}^{n_{v}} \frac{1}{n_{v}}\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}$

Remark. The moments of the $k$-point Gaussian quadrature rule $\left[\Psi_{i}\right]_{k}^{\text {gq }}$ are equal to those of $\Psi_{i}$ through degree $2 k-1$.

Theorem. The $k$-point Gaussian quadrature rule $\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}$ for $\Psi_{i}$ can be computed using $k$ iterations of the Lanczos algorithm on $\mathbf{A}$ and $\mathbf{v}_{i}$.

# Remark. The runtime for $k$ iterations of the Lanczos algorithm on $\mathbf{A}$ and $\mathbf{v}_{i}$ is $O\left(k\left(T_{\mathrm{mv}}+n\right)\right)$ 

Remark. The runtime for $k$ iterations of the Lanczos algorithm on $\mathbf{A}$ and $\mathbf{v}_{i}$ is $O\left(k\left(T_{\mathrm{mv}}+n\right)\right)$, and the required storage is $O(n)$ without reorthogonalization and $O(n k)$ with.

$$
\begin{aligned}
d_{\mathrm{W}}\left(\Phi,\left\langle\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}\right\rangle\right) & \leq d_{\mathrm{W}}\left(\Phi,\left\langle\Psi_{i}\right\rangle\right)+d_{\mathrm{W}}\left(\left\langle\Psi_{i}\right\rangle,\left\langle\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}\right\rangle\right) \\
& \leq d_{\mathrm{W}}\left(\Phi,\left\langle\Psi_{i}\right\rangle\right)+\left\langle d_{\mathrm{W}}\left(\Psi_{i},\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}\right)\right\rangle
\end{aligned}
$$

Lemma. Let $m=n \Phi(x)$. Then,

$$
\Psi_{i}(x) \sim \operatorname{Beta}\left(\frac{m}{2}, \frac{n-m}{2}\right)
$$

The uniform distribution on the unit hypersphere is unitarily invariant. Let $\mathbf{U}$ be the eigenvectors of $\mathbf{A}$. We may therefore assume

$$
\mathrm{U}^{\top} \mathbf{v}_{i} \stackrel{\mathrm{~d}}{=} \frac{\mathbf{x}}{\|\mathrm{x}\|}
$$

where $\mathbf{x} \sim \mathcal{N}(\mathbf{0}, \mathbf{I})$.

Thus, the $\left[w_{i}\right]_{j}$ have joint distribution given by,

$$
\left[w_{i}\right]_{j} \stackrel{\mathrm{~d}}{=}\left(\frac{[\mathbf{x}]_{j}}{\|\mathbf{x}\|}\right)^{2}=\frac{\left([\mathbf{x}]_{j}\right)^{2}}{\left([\mathbf{x}]_{1}\right)^{2}+\cdots+\left([\mathbf{x}]_{n}\right)^{2}}
$$

for $j=1, \ldots, n$, where $\left[\mathbf{x}_{i}\right]$ are iid standard normals.

For independent chi-square random variables $Y \sim \chi_{\alpha}^{2}$ and $Z \sim \chi_{\beta}^{2}$,

$$
\frac{Y}{Y+Z} \sim \operatorname{Beta}\left(\frac{\alpha}{2}, \frac{\beta}{2}\right) .
$$

Let $Y=\sum_{j=1}^{m}\left([\mathbf{x}]_{j}\right)^{2}$ and $Z=\sum_{j=m+1}^{n}\left([\mathbf{x}]_{j}\right)^{2}$. Then

$$
\Phi(x)=\sum_{j=1}^{m}\left[w_{i}\right]_{j}=\frac{\left([\mathbf{x}]_{1}\right)^{2}+\cdots+\left([\mathbf{x}]_{m}\right)^{2}}{\left([\mathbf{x}]_{1}\right)^{2}+\cdots+\left([\mathbf{x}]_{n}\right)^{2}}=\frac{Y}{Y+Z}
$$

Theorem. Let $X \sim \operatorname{Beta}(\alpha, \beta)$. Then $\mathbb{E}[X]=\frac{\alpha}{\alpha+\beta}$ and $X$ is $4(\alpha+\beta+1)$-sub-Gaussian.

Lemma. Suppose $X$ is $\sigma^{2}$-sub-Gaussian. Let $X_{1}, \ldots, X_{n_{v}}$ be iid samples of $X$. Then for all $t>0$,

$$
\mathbb{P}\left[\left|\left\langle X_{i}\right\rangle-\mathbb{E}[X]\right|>t\right] \leq 2 \exp \left(-\frac{n_{\vee}}{2 \sigma^{2}} t^{2}\right)
$$

Theorem. For all $t>0$,

$$
\begin{aligned}
& \max _{x} \mathbb{P}\left[\left|\Phi(x)-\left\langle\Psi_{i}(x)\right\rangle\right|>t\right] \leq 2 \exp \left(-n_{\mathrm{v}}(n+2) t^{2}\right) \\
& \mathbb{P}\left[\max _{x}\left|\Phi(x)-\left\langle\Psi_{i}(x)\right\rangle\right|>t\right] \leq 2 n \exp \left(-n_{\mathrm{v}}(n+2) t^{2}\right)
\end{aligned}
$$

Let $I[\mathbf{A}]=\lambda_{\max }-\lambda_{\min }$. Then,

$$
\begin{aligned}
d_{\mathrm{W}}\left(\Phi,\left\langle\Psi_{i}\right\rangle\right) & =\int\left|\Phi(x)-\left\langle\Psi_{i}(x)\right\rangle\right| \mathrm{d} x \\
& \leq I(\mathbf{A}) \max _{x}\left|\Phi(x)-\left\langle\Psi_{i}(x)\right\rangle\right| .
\end{aligned}
$$

Given $0<\eta<1$, if $n_{\mathrm{v}} \geq 4(n+2) t^{-2} \ln \left(2 n \eta^{-1}\right)$,

$$
\mathbb{P}\left[d_{\mathrm{W}}\left(\Phi,\left\langle\Psi_{i}\right\rangle\right)>t I[\mathbf{A}] / 2\right] \leq \eta .
$$

Proposition. If $\mu$ and $\nu$ are supported on $(a, b)$ and have equal moments through degree $s$, then

$$
d_{\mathrm{W}}(\mu, \nu) \leq 12(b-a) s^{-1}
$$

Therefore, if $k>12 t^{-1}+\frac{1}{2}$,

$$
\left\langle d_{\mathrm{W}}\left(\Psi_{i},\left[\Psi_{i}\right]_{k}^{\mathrm{gq} \mathrm{q}}\right)\right\rangle \leq d_{\mathrm{W}}\left(\Psi_{i},\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}\right) \leq t I[\mathbf{A}] / 2 .
$$

Theorem. Given $0<\eta<1$ and $t>0$, set
$n_{\mathrm{v}} \geq 4(n+2) t^{-2} \ln \left(2 n \eta^{-1}\right)$ and $k>12 t^{-1}+\frac{1}{2}$. Then, using $k n_{\mathrm{v}}$ matrix-vector products, SLQ outputs an approximation $\left\langle\left[\Psi_{i}\right]_{k}^{\mathrm{qq}}\right\rangle$ to $\Phi$ satisfying,

$$
\mathbb{P}\left[d_{\mathrm{W}}\left(\Phi,\left\langle\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}\right\rangle\right)>t I[\mathbf{A}]\right] \leq \eta .
$$

Theorem. If $\mu$ and $\nu$ are supported on $(a, b)$ and have equal moments through degree $s$, then the function $\mu-\nu$ changes sign at least $s$ times on $(a, b)$.

Since $\left[\Psi_{i}\right]_{k}^{\text {gq }}$ and $\Psi_{i}$ share $2 k-1$ moments, then $\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}-\Psi_{i}$ changes sign at least $2 k-1$ times.

The Gaussian quadrature $\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}$ is piecewise constant except at $k$ points of increase

The Gaussian quadrature $\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}$ is piecewise constant except at $k$ points of increase, and both $\Psi_{i}$ and $\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}$ are weakly increasing.

Thus, the only possible points of sign change in $\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}-\Psi_{i}$ are at the points of increase and constant regions of $\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}$

Therefore, there is a sign change in $\left[\Psi_{i}\right]_{k}^{\mathrm{gq}}-\Psi_{i}$ at every point of increase and every constant region of $[\Psi]_{k}^{\mathrm{gq}}$.


