

# Local Algorithms for Finding Densely Connected Clusters

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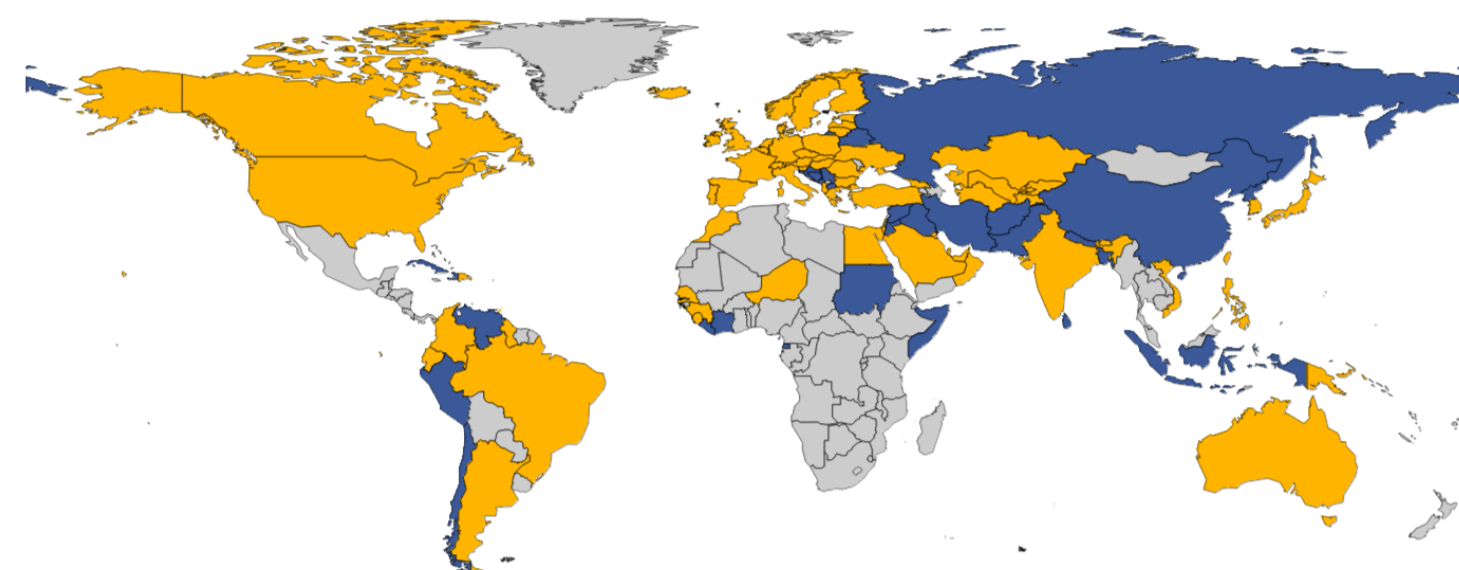
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## Problem Statement

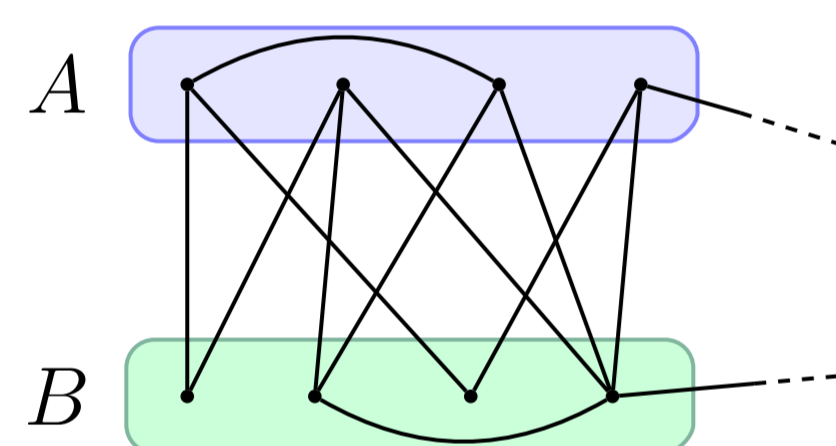
Given a graph, find two clusters  $A$  and  $B$  which are densely connected to each other and loosely connected to the rest of the graph.

**Constraint in the local setting:** the algorithm should run in time proportional to the size of returned clusters  $A$  and  $B$ .

**Application:** Interstate Dispute Graph



## Bipartiteness



- Many edges between  $A$  and  $B$
- Few edges inside  $A$  or  $B$
- Few edges to the rest of the graph

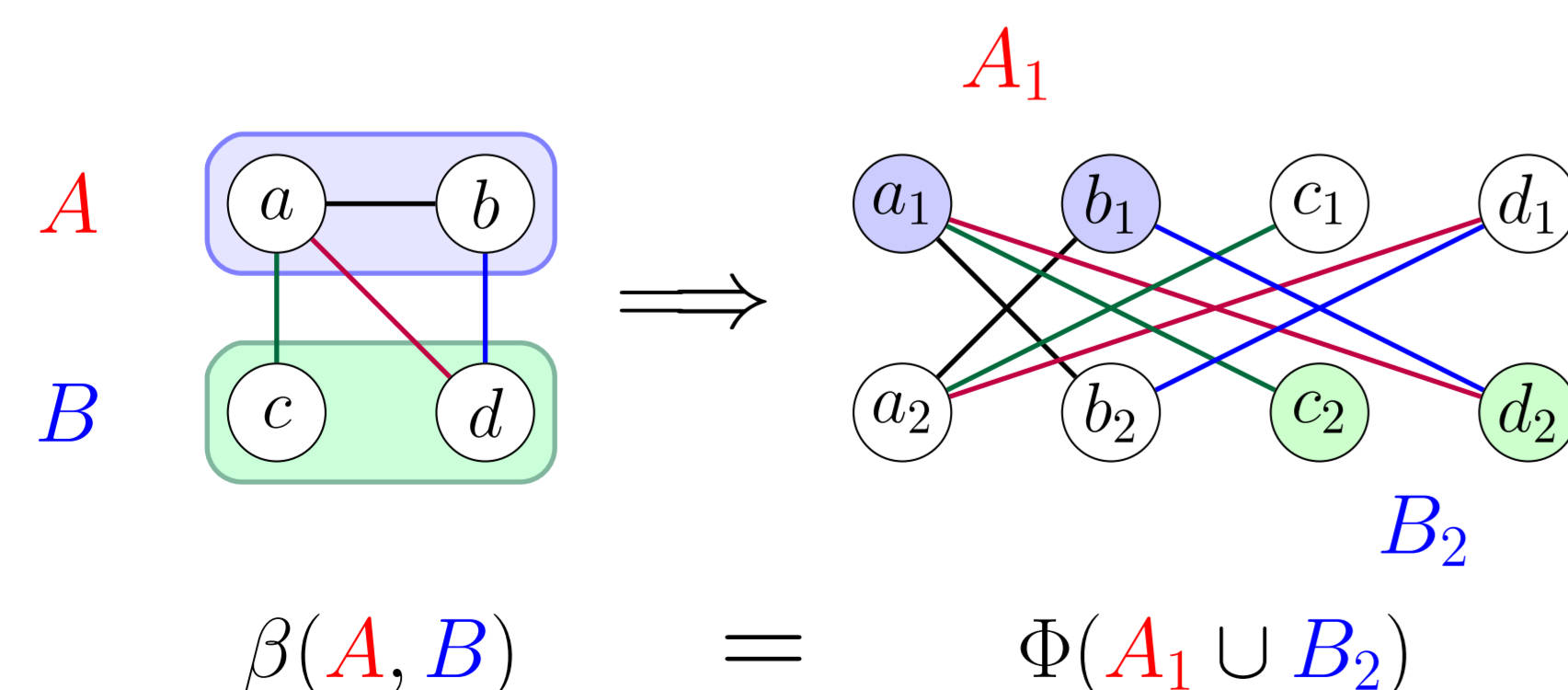
**Bipartiteness Ratio [Trevisan 2012]**

$$\beta(A, B) = 1 - \frac{2w(A, B)}{\text{vol}(A \cup B)}$$

## Reduction by Double Cover

Given a graph  $G$ , its double cover  $H$  is constructed as follows:

- every vertex  $v$  has two corresponding vertices  $v_1, v_2$ ;
- for every edge  $\{u, v\}$ , there are edges  $\{u_1, v_2\}$  and  $\{u_2, v_1\}$  in  $H$ .



$$\beta(A, B) = \Phi(A_1 \cup B_2)$$

**Lemma**

For sets  $A$  and  $B$  in the graph  $G$ , it holds that  $\beta_G(A, B) = \Phi_H(A_1 \cup B_2)$ .

## The Simplify Operator

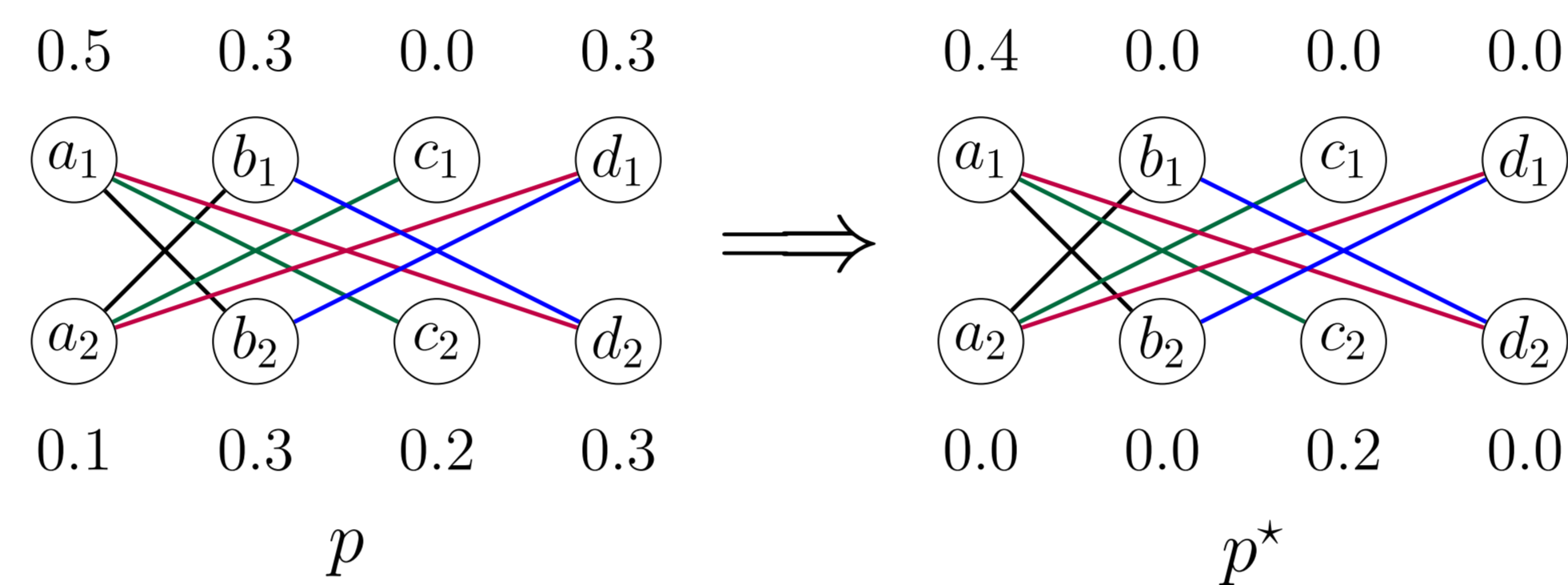
**Simplified Vector**

For any  $p \in \mathbb{R}_{\geq 0}^{2n}$ , the simplified vector  $p^*$  is defined by

$$p^*(u_1) \triangleq \max(0, p(u_1) - p(u_2)),$$

$$p^*(u_2) \triangleq \max(0, p(u_2) - p(u_1))$$

for every vertex  $u$ .



## The Algorithm for Undirected Graphs

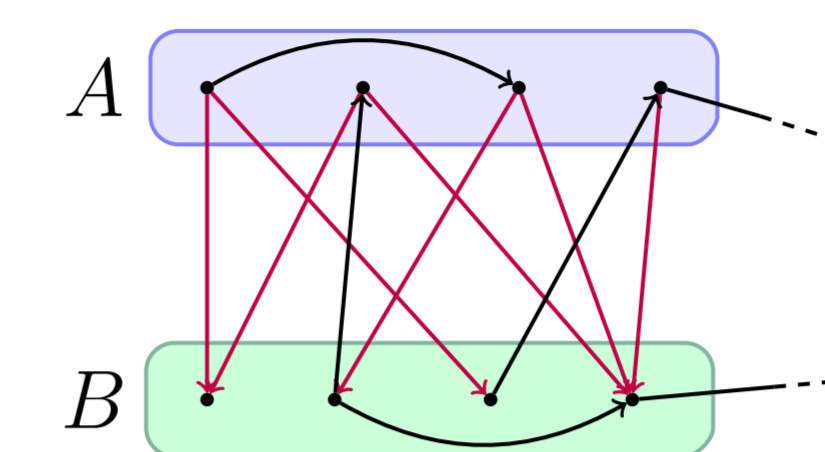
**Algorithm (Informal)**

1. Construct the double cover
2. Apply ACL-like method to compute vertex importances in the double cover
3. Compute the simplified importance vector
4. Return the vertices with high importance

**Theorem (Informal)**

Given a graph  $G$  and a starting vertex  $u$  from the target set  $A \cup B$ , there is a local algorithm which returns  $(A', B')$  with  $\beta(A', B') = O(\sqrt{\beta(A, B)})$ .

## Flow Ratio



- Many edges from  $A$  to  $B$
- Few edges inside  $A$  or  $B$
- Few edges to the rest of the graph

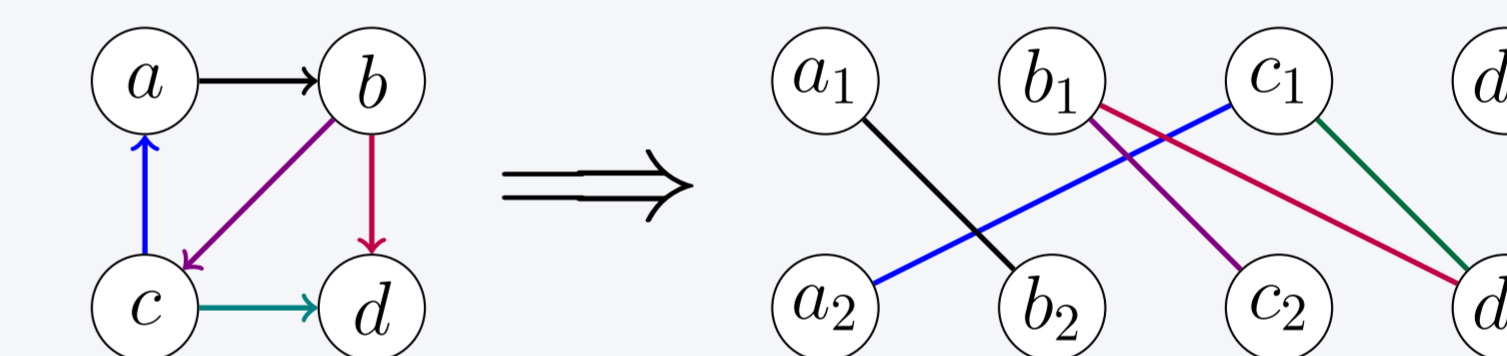
**Flow Ratio**

$$F(A, B) = 1 - \frac{2w(A, B)}{\text{vol}_{\text{out}}(A) + \text{vol}_{\text{in}}(B)}$$

## The Algorithm for Directed Graphs

Given a graph  $G$ , its semi-double cover  $H$  is constructed as follows:

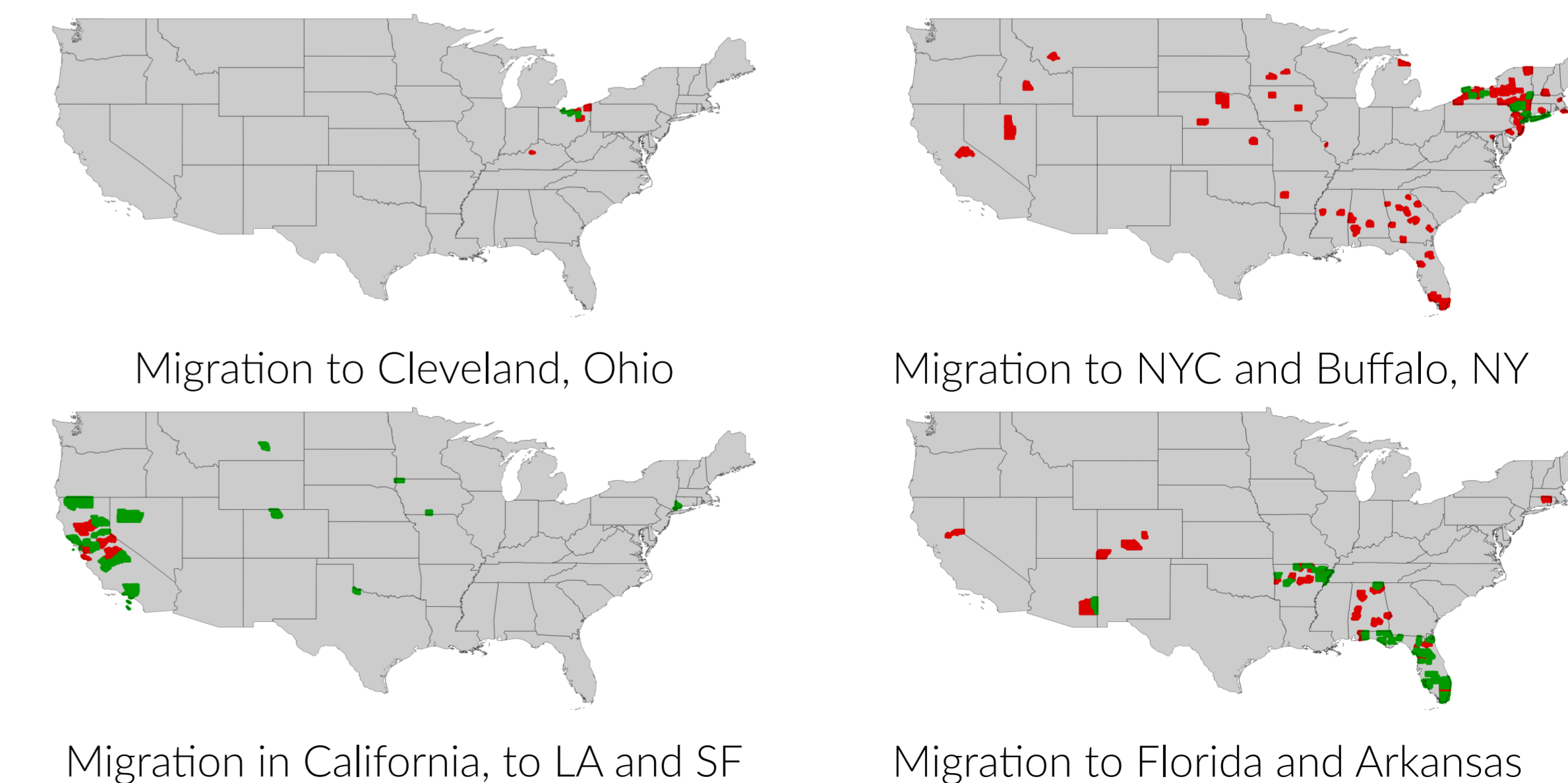
- every vertex  $v$  has two corresponding vertices  $v_1, v_2$ ;
- for every edge  $(u, v)$ , there is an edge  $\{u_1, v_2\}$  in  $H$ .



**Theorem (Informal)**

There is a local algorithm based on the semi-double cover which returns two clusters  $A$  and  $B$  with a bounded flow ratio  $F(A, B)$ .

## Experimental Results: USA Migration Graph



General migration trend from rural to urban areas

## Experimental Result: Interstate Dispute Graph

