Maximum Likelihood Estimation for Learning Populations of Parameters

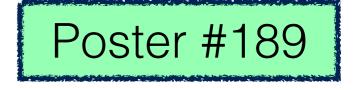
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Paul G. Allen School of CSE



joint work with Weihao Kong, Gregory Valiant, Sham Kakade

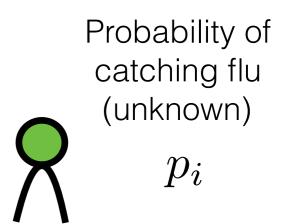


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Example: Flu data Suppose for a large random subset of the population in California, we observe whether a person caught the flu or not for last 5 years



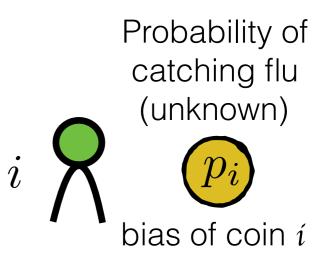
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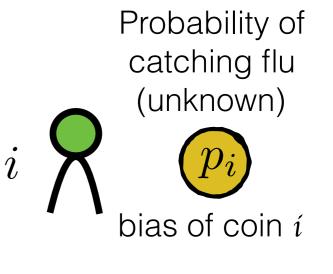


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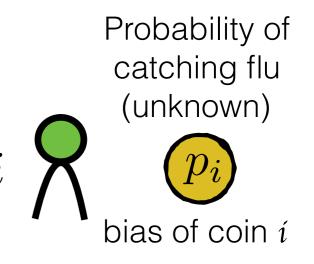


1 0) $x_i = 2$ $\widehat{p}_i = \frac{x_i}{t} = 0.4 \pm 0.45$





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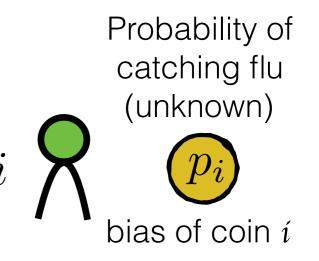


Goal: Can we learn the distribution of the biases over the population?





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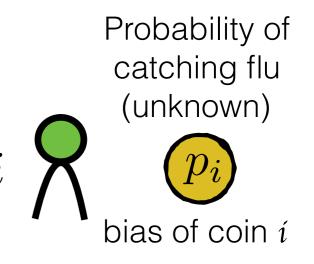
Goal: Can we learn the distribution of the biases over the population?

- Application domains: Epidemiology, Social Sciences, Psychology, Medicine, Biology
- Population size is large, often hundreds of thousands or millions
- Number of observations per individual is limited (*sparse*) prohibiting accurate estimation of parameters of interest



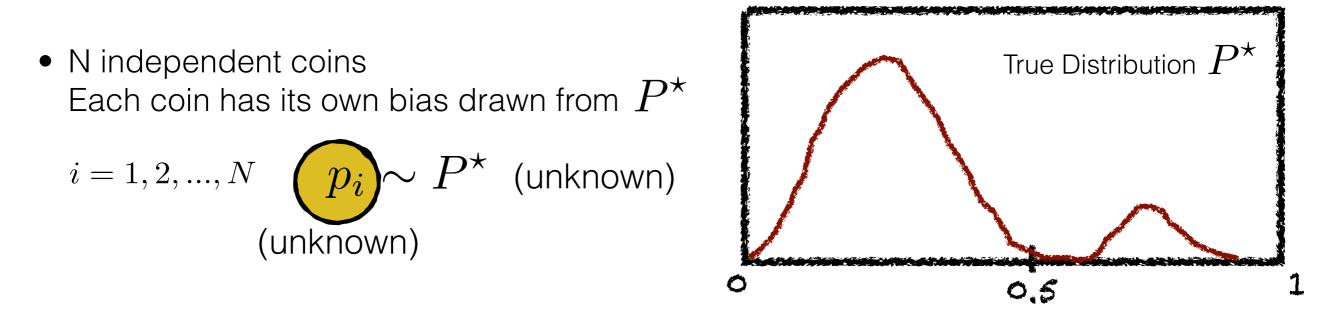


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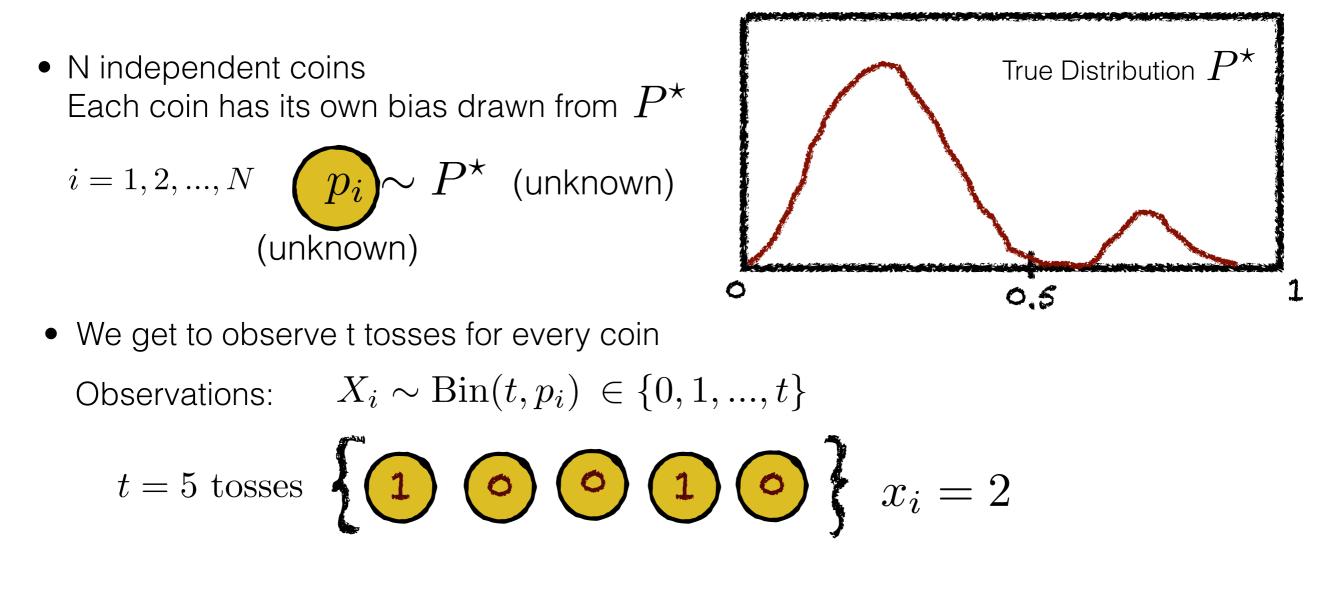


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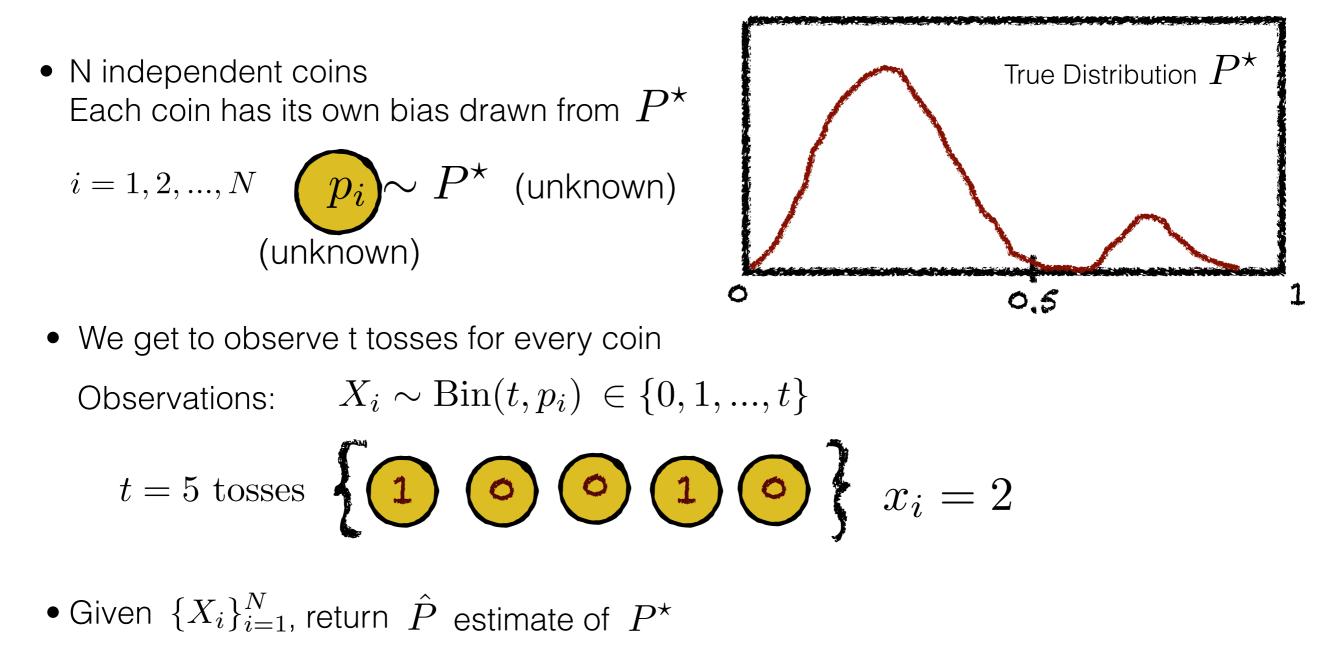
Useful for downstream analysis: Why? Testing and estimating properties of the distribution



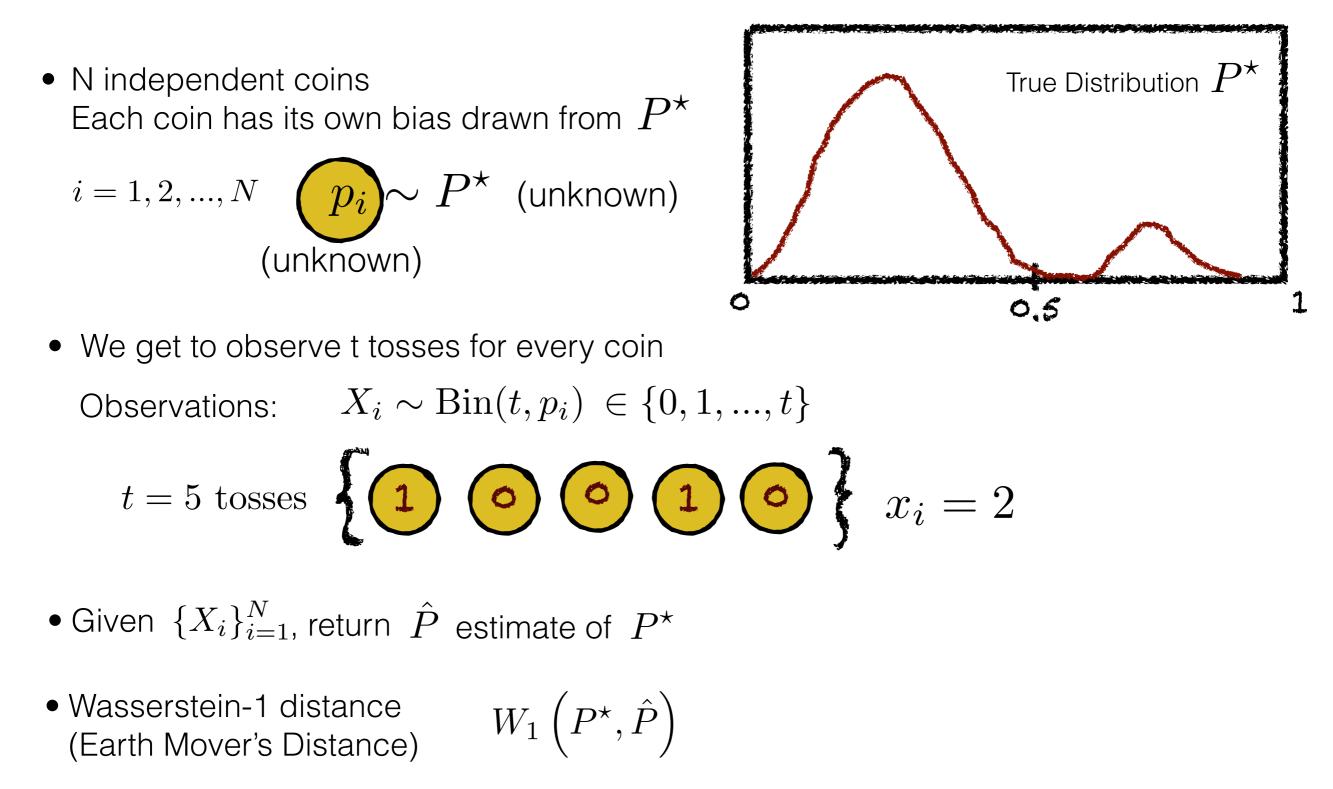














Poster #189

• Empirical plug-in estimator is bad $\hat{P}_{\text{plug-in}} = \text{histogram} \left\{ \frac{X_1}{t}, \dots, \frac{X_N}{t}, \dots, \frac{X_N}{t} \right\}$ When $t \ll N$ incurs error of $\Theta\left(\frac{1}{\sqrt{t}}\right)$ N = Number of coins t = Number of tosses per coin



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- Many recent works on estimating symmetric properties of a discrete distribution with sparse observations

Paninski 2003, Valiant and Valiant 2011, Jiao et. al. 2015, Orlitsky et. al. 2016, Acharya et. al. 2017

The setting in this work is different



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The setting in this work is different

• Tian et. al 2017 proposed a moment matching based estimator which achieves optimal error of $\mathcal{O}\left(\frac{1}{t}\right)$ when $t < c \log N$

Weakness of moment matching estimator is that it fails to obtain optimal error when $t > c \log N$ due to higher variance in larger moments



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What about Maximum Likelihood Estimator?

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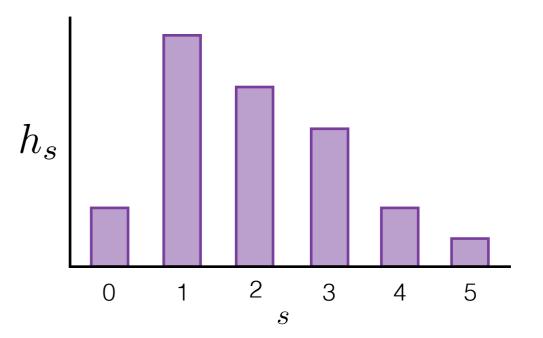


Maximum Likelihood Estimator

Sufficient statistic: Fingerprint

 $h_s = \frac{\# \text{ coins that show s heads}}{N}$ s = 0, 1, ..., t

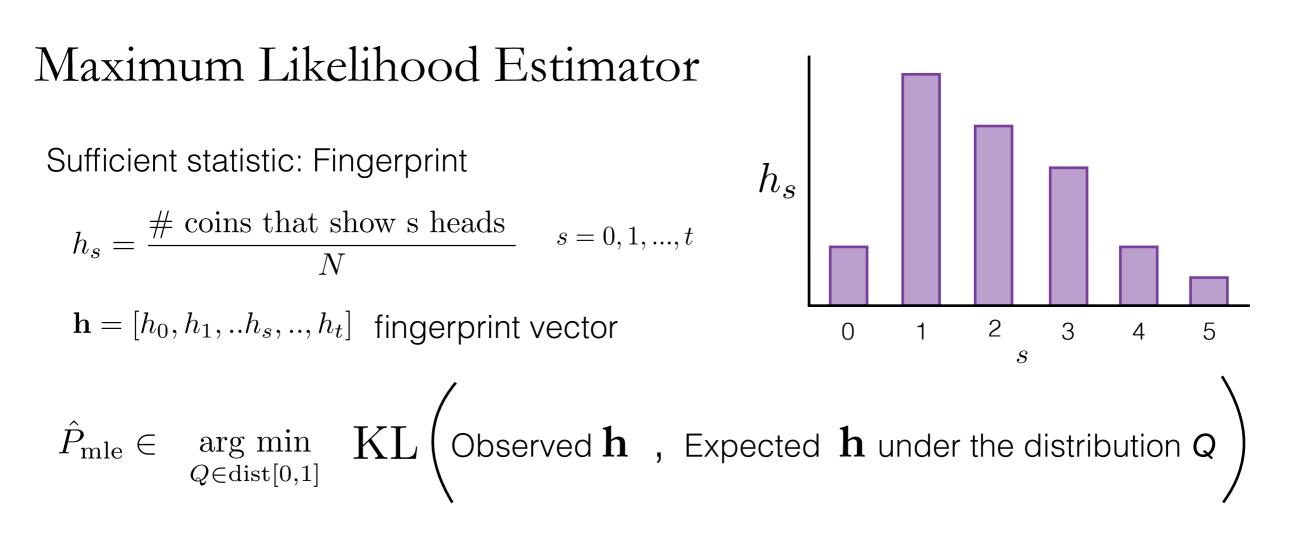
 $\mathbf{h} = [h_0, h_1, ..., h_s, ..., h_t]$ fingerprint vector





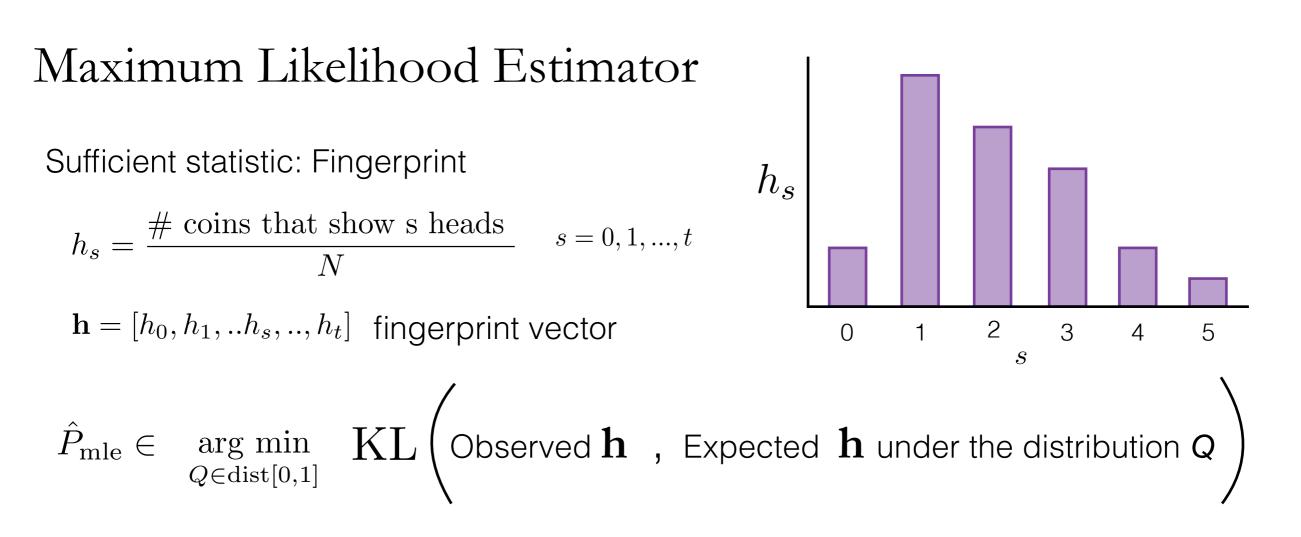
$\begin{array}{l} \text{Maximum Likelihood Estimator} \\ \text{Sufficient statistic: Fingerprint} \\ h_s = \frac{\# \text{ coins that show s heads}}{N} \quad s = 0, 1, ..., t \\ \mathbf{h} = [h_0, h_1, ..h_s, ..., h_t] \text{ fingerprint vector} \end{array} \quad h_s \left[\begin{array}{c} \mathbf{h}_s \\ \mathbf$





- NOT the empirical estimator
- Convex optimization: Efficient (polynomial time)

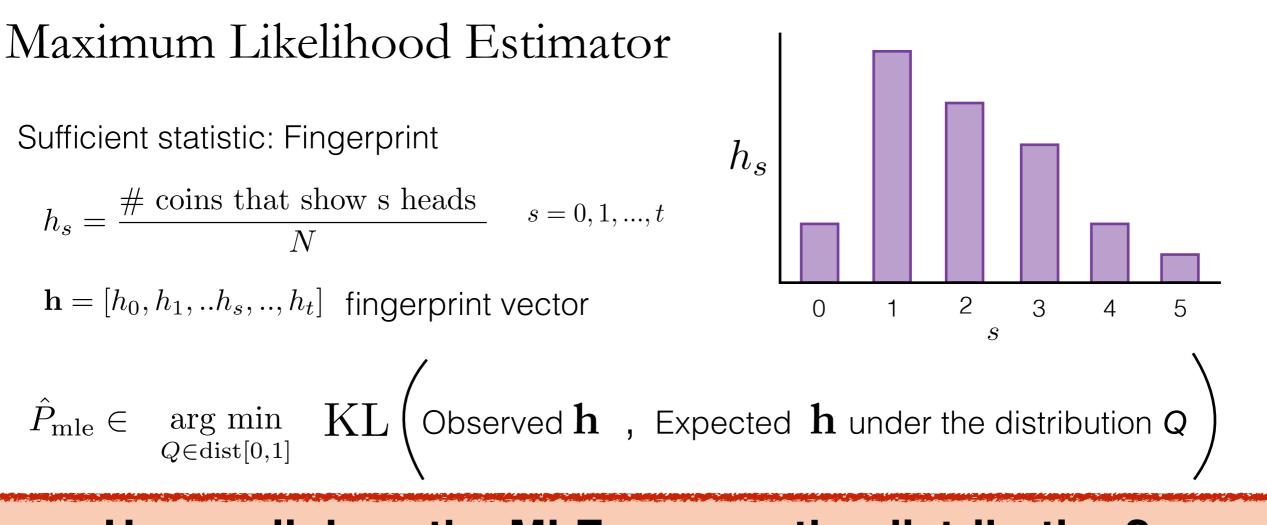




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- Convex optimization: Efficient (polynomial time)
- Proposed in late 1960's by Frederic Lord in the context of psychological testing. Several works study the geometry and identifiability and uniqueness of the solution of the MLE

Lord 1965, 1969, Turnbull 1976, Laird 1978, Lindsay 1983, Wood 1999





How well does the MLE recover the distribution?

- Convex optimization: Efficient (polynomial time)
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Main Results: MLE is Minimax Optimal in Sparse Regime

Non-asymptotic guarantees

Theorem 1

The MLE achieves following error bounds: w. p. $\geq 1 - \delta$

• Small Sample Regime:

$$W_1\left(P^{\star}, \hat{P}_{\text{mle}}\right) = \mathcal{O}_{\delta}\left(\frac{1}{t}\right) \text{ when } t < c \log N$$

• Medium Sample Regime:

$$W_1\left(P^\star, \hat{P}_{\rm mle}\right) = \mathcal{O}_{\delta}\left(\frac{1}{\sqrt{t \, \log N}}\right) \text{ when } c \log N \le t \le N^{2/9-\epsilon}$$

$$N = {egin{array}{c} {
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Theorem 2

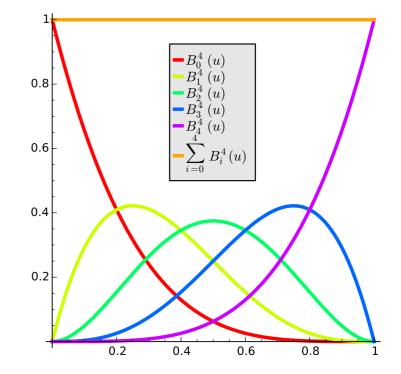
• Matching Minimax Lower Bounds

$$\inf_{f} \sup_{P} \operatorname{E} \left[W_1(P, f(\mathbf{X})) \right] > \Omega\left(\frac{1}{t}\right) \vee \Omega\left(\frac{1}{\sqrt{t \log N}}\right)$$

Novel Proof: Polynomial Approximations

New bounds on coefficients of Bernstein polynomials approximating Lipschitz-1 functions on [0, 1]

$$\widehat{f}_t(x) = \sum_{j=0}^t b_j \binom{t}{j} x^j (1-x)^{t-j}$$



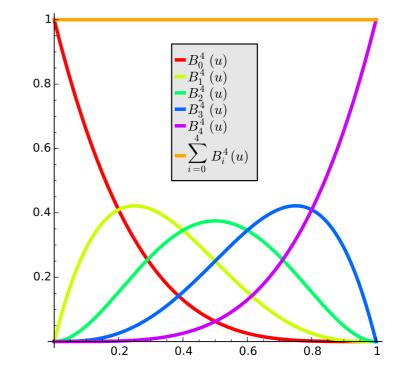
Bernstein polynomials



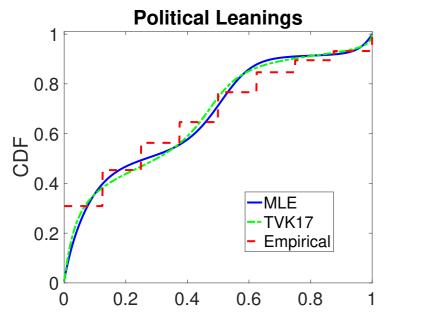
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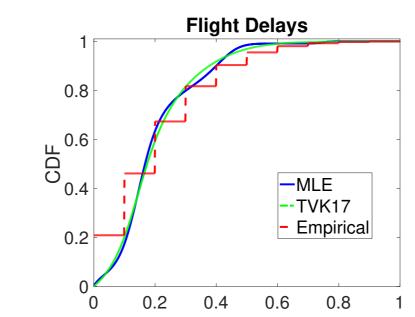
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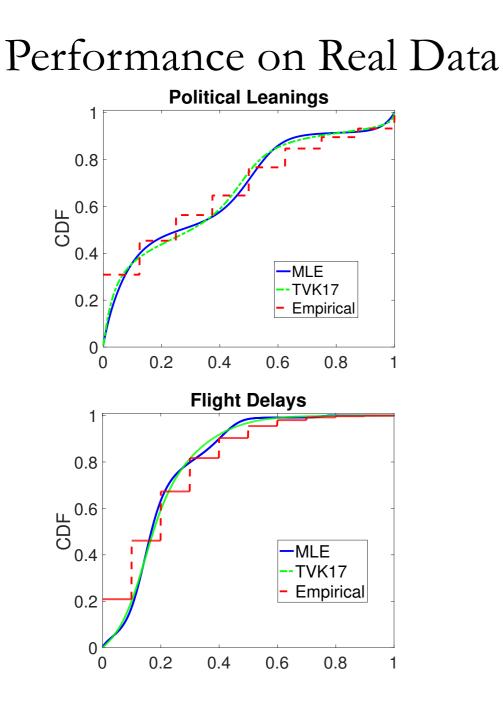
Performance on Real Data





Summary

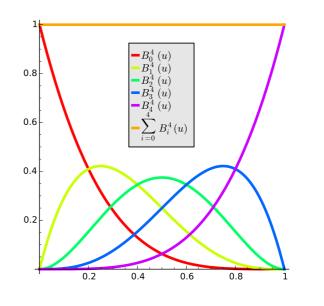
Learning distribution of parameters over a population with sparse observations per individual



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MLE is Minimax Optimal even with sparse observations!

Novel proof: new bounds on coefficients of Bernstein polynomials approximating Lipschitz-1 functions



Poster #189