# Approximating Orthogonal Matrices with Effective Givens Factorization

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joint work with Joan Bruna (NYU)

Poster #164

### **Givens Factorization of Orthogonal Matrices**

$$G^{T}(i,j,\alpha) = \begin{bmatrix} 1 & \cdots & 0 & \cdots & 0 & \cdots & 0 \\ \vdots & \ddots & \vdots & & \vdots & & \vdots \\ 0 & \cdots & \cos(\alpha) & \cdots & -\sin(\alpha) & \cdots & 0 \\ \vdots & \vdots & \ddots & \vdots & & \vdots \\ 0 & \cdots & \sin(\alpha) & \cdots & \cos(\alpha) & \cdots & 0 \\ \vdots & \vdots & & \vdots & \ddots & \vdots \\ 0 & \cdots & 0 & \cdots & 0 & \cdots & 1 \end{bmatrix}$$

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**Exact Givens Factorization** 

$$U = G_1 \dots G_N$$
  $N = \frac{d(d-1)}{2}$ 

**Approximate Givens Factorization** 

$$U \approx G_1 \dots G_N$$
  $N \ll \frac{d(d-1)}{2}$   
computationally hard problem

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**Our Questions in this Context** 

 Which orthogonal matrices can be effectively approximated? (not all of them) **Approximate Givens Factorization** 

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**Our Questions in this Context** 

- Which orthogonal matrices can be effectively approximated? (not all of them)
- 2. Which principles are behind effective approximation algorithms? (sparsity-inducing algorithms)

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Once computed, applied many times

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**Unitary Basis Transform** 

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#### **Advantageous Setting**

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**Unitary Basis Transform** 

$$\mathsf{FFT} \colon \mathcal{O}\left(d^2\right) \to \mathcal{O}\left(d\log(d)\right)$$

Application: Graph Fourier Transform

#### Which Matrices can be Effectively Approximated?

#### Theorem

Let  $\epsilon > 0$ . If  $N = o(d^2/\log(d))$ , then as  $d \to \infty$ ,

$$\mu\left(\left\{\left.U\in U(d)\right|\inf_{G_1\ldots G_N}\|U-\prod_n G_n\|_2\leq\epsilon
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where  $\mu$  is the Haar measure over U(d).

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- proof is based on an  $\epsilon$ -covering argument
- suggests computational-to-statistical gap together with experimental results (details at poster)

### K-planted Distribution over SO(d)

Sample  $U = G_1 \dots G_K$ 

- choose subspace  $(i_k, j_k)$  uniformly with replacement
- choose rotation angle  $\alpha_k \in [0, 2\pi)$  uniformly

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*K*-planted matrices quickly become dense

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**Approximation criterion** 

$$\left\| U - \hat{U} \right\|_{F, \operatorname{sym}} \coloneqq \min_{P \in \mathcal{P}_d} \left\| U - \hat{U}P \right\|_F$$

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$$f(U) \coloneqq d^{-1} \|U\|_1 = d^{-1} \sum_{i,j=1}^d |U_{ij}|$$

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- *Non-convex* greedy step
- global optimum in  $\mathcal{O}(d^2)$  amortized time complexity

## Thank you

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https://github.com/tfrerix/givens-factorization