#### Boosting Frank-Wolfe by Chasing Gradients

#### Cyrille W. Combettes

with Sebastian Pokutta

School of Industrial and Systems Engineering Georgia Institute of Technology Atlanta, GA, USA

37<sup>th</sup> International Conference on Machine Learning July 12–18, 2020





- **2** The Frank-Wolfe algorithm
- **3** Boosting Frank-Wolfe
- **4** Computational experiments

## Let $\mathcal{H}$ be a Euclidean space (e.g., $\mathbb{R}^n$ or $\mathbb{R}^{m \times n}$ ) and consider min f(x)s.t. $x \in C$

where

- $f: \mathcal{H} \to \mathbb{R}$  is a smooth convex function
- $\mathcal{C} \subset \mathcal{H}$  is a compact convex set,  $\mathcal{C} = \mathsf{conv}(\mathcal{V})$

# Let $\mathcal H$ be a Euclidean space (e.g., $\mathbb R^n$ or $\mathbb R^{m\times n})$ and consider

min f(x)s.t.  $x \in C$ 

where

- $f: \mathcal{H} \to \mathbb{R}$  is a smooth convex function
- $\mathcal{C} \subset \mathcal{H}$  is a compact convex set,  $\mathcal{C} = \operatorname{conv}(\mathcal{V})$

#### Example

• Sparse logistic regression

$$\min_{x \in \mathbb{R}^n} \frac{1}{m} \sum_{i=1}^m \ln(1 + \exp(-y_i a_i^\top x))$$
  
s.t.  $\|x\|_1 \leqslant \tau$ 

• Low-rank matrix completion

$$\min_{\substack{X \in \mathbb{R}^{m \times n}}} \frac{1}{2|\mathcal{I}|} \sum_{(i,j) \in \mathcal{I}} (Y_{i,j} - X_{i,j})^2$$
s.t.  $\|X\|_{\mathsf{nuc}} \leqslant \tau$ 

• A natural approach is to use any efficient method and add projections back onto C to ensure feasibility

• A natural approach is to use any efficient method and add projections back onto C to ensure feasibility



- A natural approach is to use any efficient method and add projections back onto C to ensure feasibility
- However, in many situations projections onto  $\mathcal C$  are very expensive

- A natural approach is to use any efficient method and add projections back onto C to ensure feasibility
- However, in many situations projections onto  $\mathcal C$  are very expensive
- This is an issue with the method of projections, not necessarily with the geometry of C: linear minimizations over C can still be relatively cheap

- A natural approach is to use any efficient method and add projections back onto  ${\cal C}$  to ensure feasibility
- However, in many situations projections onto  ${\mathcal C}$  are very expensive
- This is an issue with the method of projections, not necessarily with the geometry of C: linear minimizations over C can still be relatively cheap

Feasible region $\mathcal{C}$	Linear minimization	Projection
$\ell_1/\ell_2/\ell_\infty$ -ball	$\mathcal{O}(n)$	$\mathcal{O}(n)$
$\ell_{p}$ -ball, $p \in ]1,\infty[\setminus\{2\}]$	$\mathcal{O}(n)$	N/A
Nuclear norm-ball	$\mathcal{O}(nnz)$	$\mathcal{O}(mn\min\{m,n\})$
Flow polytope	$\mathcal{O}(n)$	$\mathcal{O}(n^{3.5})$
Birkhoff polytope	$\mathcal{O}(n^3)$	N/A
Matroid polytope	$\mathcal{O}(n \ln(n))$	$\mathcal{O}(poly(n))$

- A natural approach is to use any efficient method and add projections back onto  ${\cal C}$  to ensure feasibility
- However, in many situations projections onto  $\mathcal C$  are very expensive
- This is an issue with the method of projections, not necessarily with the geometry of C: linear minimizations over C can still be relatively cheap

Feasible region $\mathcal{C}$	Linear minimization	Projection
$\ell_1/\ell_2/\ell_\infty$ -ball	$\mathcal{O}(n)$	$\mathcal{O}(n)$
$\ell_p$ -ball, $p \in ]1,\infty[\setminus\{2\}]$	$\mathcal{O}(n)$	N/A
Nuclear norm-ball	$\mathcal{O}(nnz)$	$\mathcal{O}(mn\min\{m,n\})$
Flow polytope	$\mathcal{O}(n)$	$\mathcal{O}(n^{3.5})$
Birkhoff polytope	$\mathcal{O}(n^3)$	N/A
Matroid polytope	$\mathcal{O}(n \ln(n))$	$\mathcal{O}(poly(n))$

- A natural approach is to use any efficient method and add projections back onto C to ensure feasibility
- However, in many situations projections onto  $\mathcal C$  are very expensive
- This is an issue with the method of projections, not necessarily with the geometry of C: linear minimizations over C can still be relatively cheap

Feasible region $\mathcal{C}$	Linear minimization	Projection
$\ell_1/\ell_2/\ell_\infty$ -ball	$\mathcal{O}(n)$	$\mathcal{O}(n)$
$\ell_p$ -ball, $p \in ]1, \infty[\setminus \{2\}]$	$\mathcal{O}(n)$	N/A
Nuclear norm-ball	$\mathcal{O}(nnz)$	$\mathcal{O}(mn\min\{m,n\})$
Flow polytope	$\mathcal{O}(n)$	$\mathcal{O}(n^{3.5})$
Birkhoff polytope	$\mathcal{O}(n^3)$	N/A
Matroid polytope	$\mathcal{O}(n\ln(n))$	$\mathcal{O}(poly(n))$

- A natural approach is to use any efficient method and add projections back onto C to ensure feasibility
- However, in many situations projections onto  $\mathcal C$  are very expensive
- This is an issue with the method of projections, not necessarily with the geometry of C: linear minimizations over C can still be relatively cheap

Feasible region $\mathcal{C}$	Linear minimization	Projection
$\ell_1/\ell_2/\ell_\infty$ -ball	$\mathcal{O}(n)$	$\mathcal{O}(n)$
$\ell_{p}$ -ball, $p \in ]1,\infty[\setminus\{2\}]$	$\mathcal{O}(n)$	N/A
Nuclear norm-ball	$\mathcal{O}(nnz)$	$\mathcal{O}(mn\min\{m,n\})$
Flow polytope	$\mathcal{O}(n)$	$\mathcal{O}(n^{3.5})$
Birkhoff polytope	$\mathcal{O}(n^3)$	N/A
Matroid polytope	$\mathcal{O}(n \ln(n))$	$\mathcal{O}(poly(n))$

- A natural approach is to use any efficient method and add projections back onto  ${\cal C}$  to ensure feasibility
- However, in many situations projections onto  ${\mathcal C}$  are very expensive
- This is an issue with the method of projections, not necessarily with the geometry of C: linear minimizations over C can still be relatively cheap

Feasible region $\mathcal{C}$	Linear minimization	Projection
$\ell_1/\ell_2/\ell_\infty$ -ball	$\mathcal{O}(n)$	$\mathcal{O}(n)$
$\ell_p$ -ball, $p \in ]1,\infty[\setminus\{2\}]$	$\mathcal{O}(n)$	N/A
Nuclear norm-ball	$\mathcal{O}(nnz)$	$\mathcal{O}(mn\min\{m,n\})$
Flow polytope	$\mathcal{O}(n)$	$\mathcal{O}(n^{3.5})$
Birkhoff polytope	$\mathcal{O}(n^3)$	N/A
Matroid polytope	$\mathcal{O}(n \ln(n))$	$\mathcal{O}(poly(n))$

N/A: no closed-form exists and solution must be computed via nontrivial optimization

• Can we avoid projections?

The Frank-Wolfe algorithm (Frank & Wolfe, 1956) a.k.a. conditional gradient algorithm (Levitin & Polyak, 1966):

AlgorithmFrank-Wolfe (FW)Input:  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ .1: for t = 0 to T - 1 do2:  $v_t \leftarrow \underset{v \in \mathcal{V}}{\operatorname{arg\,min}} \langle \nabla f(x_t), v \rangle$ 3:  $x_{t+1} \leftarrow x_t + \gamma_t (v_t - x_t)$ 



The Frank-Wolfe algorithm (Frank & Wolfe, 1956) a.k.a. conditional gradient algorithm (Levitin & Polyak, 1966):

AlgorithmFrank-Wolfe (FW)Input:  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ .1: for t = 0 to T - 1 do2:  $v_t \leftarrow \underset{v \in \mathcal{V}}{\operatorname{arg\,min}} \langle \nabla f(x_t), v \rangle$ 3:  $x_{t+1} \leftarrow x_t + \gamma_t (v_t - x_t)$ 



The Frank-Wolfe algorithm (Frank & Wolfe, 1956) a.k.a. conditional gradient algorithm (Levitin & Polyak, 1966):

AlgorithmFrank-Wolfe (FW)Input:  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ .1: for t = 0 to T - 1 do2:  $v_t \leftarrow \operatorname*{arg\,min}_{v \in \mathcal{V}} \langle \nabla f(x_t), v \rangle$ 3:  $x_{t+1} \leftarrow x_t + \gamma_t (v_t - x_t)$ 



The Frank-Wolfe algorithm (Frank & Wolfe, 1956) a.k.a. conditional gradient algorithm (Levitin & Polyak, 1966):

AlgorithmFrank-Wolfe (FW)Input:  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ .1: for t = 0 to T - 1 do2:  $v_t \leftarrow \operatorname*{arg\,min}_{v \in \mathcal{V}} \langle \nabla f(x_t), v \rangle$ 3:  $x_{t+1} \leftarrow x_t + \gamma_t (v_t - x_t)$ 



The Frank-Wolfe algorithm (Frank & Wolfe, 1956) a.k.a. conditional gradient algorithm (Levitin & Polyak, 1966):

AlgorithmFrank-Wolfe (FW)Input:  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ .1: for t = 0 to T - 1 do2:  $v_t \leftarrow \arg\min(\nabla f(x_t), v)$ 3:  $x_{t+1} \leftarrow x_t + \gamma_t(v_t - x_t)$ 



The Frank-Wolfe algorithm (Frank & Wolfe, 1956) a.k.a. conditional gradient algorithm (Levitin & Polyak, 1966):

AlgorithmFrank-Wolfe (FW)Input:  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ .1: for t = 0 to T - 1 do2:  $v_t \leftarrow \arg\min_{v \in \mathcal{V}} \langle \nabla f(x_t), v \rangle$ 3:  $x_{t+1} \leftarrow x_t + \gamma_t(v_t - x_t)$ 



The Frank-Wolfe algorithm (Frank & Wolfe, 1956) a.k.a. conditional gradient algorithm (Levitin & Polyak, 1966):

AlgorithmFrank-Wolfe (FW)Input:  $x_0 \in C, \ \gamma_t \in [0, 1].$ 1: for t = 0 to T - 1 do2:  $v_t \leftarrow \arg\min_{v \in V} \langle \nabla f(x_t), v \rangle$ 3:  $x_{t+1} \leftarrow x_t + \gamma_t (v_t - x_t)$ 



•  $x_{t+1}$  is obtained by convex combination of  $x_t \in C$  and  $v_t \in C$ , thus  $x_{t+1} \in C$ 

The Frank-Wolfe algorithm (Frank & Wolfe, 1956) a.k.a. conditional gradient algorithm (Levitin & Polyak, 1966):

AlgorithmFrank-Wolfe (FW)Input:  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ .1: for t = 0 to T - 1 do2:  $v_t \leftarrow \arg\min_{v \in V} \langle \nabla f(x_t), v \rangle$ 3:  $x_{t+1} \leftarrow x_t + \gamma_t (v_t - x_t)$ 



- $x_{t+1}$  is obtained by convex combination of  $x_t \in C$  and  $v_t \in C$ , thus  $x_{t+1} \in C$
- FW uses linear minimizations (the "FW oracle") instead of projections

The Frank-Wolfe algorithm (Frank & Wolfe, 1956) a.k.a. conditional gradient algorithm (Levitin & Polyak, 1966):

AlgorithmFrank-Wolfe (FW)Input:  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ .1: for t = 0 to T - 1 do2:  $v_t \leftarrow \arg\min_{v \in V} \langle \nabla f(x_t), v \rangle$ 3:  $x_{t+1} \leftarrow x_t + \gamma_t (v_t - x_t)$ 



- $x_{t+1}$  is obtained by convex combination of  $x_t \in C$  and  $v_t \in C$ , thus  $x_{t+1} \in C$
- FW uses linear minimizations (the "FW oracle") instead of projections
- FW = pick a vertex (using gradient information) and move in that direction

The Frank-Wolfe algorithm (Frank & Wolfe, 1956) a.k.a. conditional gradient algorithm (Levitin & Polyak, 1966):

**Algorithm** Frank-Wolfe (FW) **Input:**  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ . 1: for t = 0 to T - 1 do 2:  $v_t \leftarrow \underset{v \in \mathcal{V}}{\arg\min} \langle \nabla f(x_t), v \rangle$  $x_{t+1} \leftarrow x_t + \gamma_t (v_t - x_t)$ 





- $x_{t+1}$  is obtained by convex combination of  $x_t \in C$  and  $v_t \in C$ , thus  $x_{t+1} \in C$
- FW uses linear minimizations (the "FW oracle") instead of projections
- FW = pick a vertex (using gradient information) and move in that direction
- Successfully applied to: traffic assignment, computer vision, optimal transport, adversarial learning, etc.

#### Theorem (Levitin & Polyak, 1966; Jaggi, 2013)

Let  $C \subset \mathcal{H}$  be a compact convex set with diameter D and  $f : \mathcal{H} \to \mathbb{R}$  be a L-smooth convex function, and let  $x_0 \in \arg\min_{v \in \mathcal{V}} \langle \nabla f(y), v \rangle$  for some  $y \in C$ . If  $\gamma_t = \frac{2}{t+2}$  (default) or  $\gamma_t = \min\left\{\frac{\langle \nabla f(x_t), x_t - v_t \rangle}{L \|x_t - v_t\|^2}, 1\right\}$  ("short step"), then

$$f(x_t) - \min_{\mathcal{C}} f \leqslant \frac{4LD^2}{t+2}$$

#### Theorem (Levitin & Polyak, 1966; Jaggi, 2013)

Let  $C \subset \mathcal{H}$  be a compact convex set with diameter D and  $f : \mathcal{H} \to \mathbb{R}$  be a *L*-smooth convex function, and let  $x_0 \in \arg\min_{v \in \mathcal{V}} \langle \nabla f(y), v \rangle$  for some  $y \in C$ . If  $\gamma_t = \frac{2}{t+2}$  (default) or  $\gamma_t = \min\left\{\frac{\langle \nabla f(x_t), x_t - v_t \rangle}{L \|x_t - v_t\|^2}, 1\right\}$  ("short step"), then

$$f(x_t) - \min_{\mathcal{C}} f \leqslant \frac{4LD^2}{t+2}$$

• The convergence rate cannot be improved (Canon & Cullum, 1968; Jaggi, 2013; Lan, 2013)

#### Theorem (Levitin & Polyak, 1966; Jaggi, 2013)

Let  $C \subset \mathcal{H}$  be a compact convex set with diameter D and  $f : \mathcal{H} \to \mathbb{R}$  be a *L*-smooth convex function, and let  $x_0 \in \arg\min_{v \in \mathcal{V}} \langle \nabla f(y), v \rangle$  for some  $y \in C$ . If  $\gamma_t = \frac{2}{t+2}$  (default) or  $\gamma_t = \min \left\{ \frac{\langle \nabla f(x_t), x_t - v_t \rangle}{L \|x_t - v_t\|^2}, 1 \right\}$  ("short step"), then

$$f(x_t) - \min_{\mathcal{C}} f \leqslant \frac{4LD^2}{t+2}$$

- The convergence rate cannot be improved (Canon & Cullum, 1968; Jaggi, 2013; Lan, 2013)
- Why?

Consider the simple problem

$$\min \frac{1}{2} \|x\|_2^2$$
 s.t.  $x \in \operatorname{conv}\left(\begin{pmatrix}0\\1\end{pmatrix}, \begin{pmatrix}-1\\0\end{pmatrix}, \begin{pmatrix}1\\0\end{pmatrix}\right)$ 

and 
$$x^* = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$$



Consider the simple problem

$$\begin{split} &\min \frac{1}{2} \|x\|_2^2 \\ &\text{s.t. } x \in \operatorname{conv} \left( \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \begin{pmatrix} -1 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \end{pmatrix} \right) \end{split}$$

and  $x^* = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$ 

• Let 
$$x_0 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$







• Let 
$$x_0 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$





• Let 
$$x_0 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

Consider the simple problem

 $\begin{array}{l} \min \frac{1}{2} \|x\|_2^2 \\ \text{s.t. } x \in \operatorname{conv} \left( \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \begin{pmatrix} -1 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \end{pmatrix} \right) \\ \text{and } x^* = \begin{pmatrix} 0 \\ 0 \end{pmatrix} \end{array}$ 

x<sub>1</sub>

• Let 
$$x_0 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$



• Let 
$$x_0 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$



Consider the simple problem

 $\min \frac{1}{2} \|x\|_{2}^{2}$ s.t.  $x \in \operatorname{conv}\left(\begin{pmatrix}0\\1\end{pmatrix}, \begin{pmatrix}-1\\0\end{pmatrix}, \begin{pmatrix}1\\0\end{pmatrix}\right)$ 

and  $x^* = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$ 

• Let 
$$x_0 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$





• Let 
$$x_0 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$



Consider the simple problem

 $\begin{array}{l} \min \frac{1}{2} \|x\|_2^2 \\ \text{s.t. } x \in \operatorname{conv} \left( \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \begin{pmatrix} -1 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \end{pmatrix} \right) \\ \text{and } x^* = \begin{pmatrix} 0 \\ 0 \end{pmatrix} \end{array}$ 

• Let 
$$x_0 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$



Consider the simple problem

 $\min \frac{1}{2} \|x\|_2^2$ s.t.  $x \in \operatorname{conv}\left(\begin{pmatrix}0\\1\end{pmatrix}, \begin{pmatrix}-1\\0\end{pmatrix}, \begin{pmatrix}1\\0\end{pmatrix}\right)$ and  $x^* = \begin{pmatrix}0\\0\end{pmatrix}$ 

• Let 
$$x_0 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$
### The Frank-Wolfe algorithm





• Let 
$$x_0 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

• FW tries to reach x\* by moving towards vertices

## The Frank-Wolfe algorithm

Consider the simple problem

 $\begin{array}{l} \min \, \frac{1}{2} \|x\|_2^2 \\ \text{s.t. } x \in \operatorname{conv} \left( \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \begin{pmatrix} -1 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \end{pmatrix} \right) \end{array}$ 

and  $x^* = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$ 

• Let 
$$x_0 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$

- FW tries to reach x\* by moving towards vertices
- This yields an inefficient zig-zagging trajectory



• Away-Step Frank-Wolfe (AFW) (Wolfe, 1970; Lacoste-Julien & Jaggi, 2015): enhances FW by allowing to move away from vertices

• Away-Step Frank-Wolfe (AFW) (Wolfe, 1970; Lacoste-Julien & Jaggi, 2015): enhances FW by allowing to move away from vertices



• Away-Step Frank-Wolfe (AFW) (Wolfe, 1970; Lacoste-Julien & Jaggi, 2015): enhances FW by allowing to move away from vertices



 Decomposition-Invariant Pairwise Conditional Gradient (DICG) (Garber & Meshi, 2016): memory-free variant of AFW

• Away-Step Frank-Wolfe (AFW) (Wolfe, 1970; Lacoste-Julien & Jaggi, 2015): enhances FW by allowing to move away from vertices



- Decomposition-Invariant Pairwise Conditional Gradient (DICG) (Garber & Meshi, 2016): memory-free variant of AFW
- Blended Conditional Gradients (BCG) (Braun et al., 2019): blends FCFW and FW

• Can we speed up FW in a simple way?

- Can we speed up FW in a simple way?
- Rule of thumb in optimization: follow the steepest direction

- Can we speed up FW in a simple way?
- Rule of thumb in optimization: follow the steepest direction

Idea:

• Speed up FW by moving in a direction better aligned with  $-\nabla f(x_t)$ 

- Can we speed up FW in a simple way?
- Rule of thumb in optimization: follow the steepest direction

Idea:

- Speed up FW by moving in a direction better aligned with  $-\nabla f(x_t)$
- Build this direction by using  ${\mathcal V}$  to maintain the projection-free property



• 
$$v_0 \in \arg \max_{v \in \mathcal{V}} \langle -\nabla f(x_t), v \rangle$$
  
 $\lambda_0 u_0 = \frac{\langle -\nabla f(x_t), v_0 - x_t \rangle}{\|v_0 - x_t\|^2} (v_0 - x_t)$   
 $r_1 = -\nabla f(x_t) - \lambda_0 u_0$ 



• 
$$v_0 \in \arg \max_{v \in \mathcal{V}} \langle -\nabla f(x_t), v \rangle$$
  
 $\lambda_0 u_0 = \frac{\langle -\nabla f(x_t), v_0 - x_t \rangle}{\|v_0 - x_t\|^2} (v_0 - x_t)$   
 $r_1 = -\nabla f(x_t) - \lambda_0 u_0$ 



 How can we build a direction better aligned with −∇f(x<sub>t</sub>) and that allows to update x<sub>t+1</sub> without projection?

• 
$$v_0 \in \arg \max_{v \in \mathcal{V}} \langle -\nabla f(x_t), v \rangle$$
  
 $\lambda_0 u_0 = \frac{\langle -\nabla f(x_t), v_0 - x_t \rangle}{\|v_0 - x_t\|^2} (v_0 - x_t)$   
 $r_1 = -\nabla f(x_t) - \lambda_0 u_0$ 

•  $v_1 \in \arg \max_{v \in \mathcal{V}} \langle r_1, v \rangle$   $\lambda_1 u_1 = \frac{\langle r_1, v_1 - x_t \rangle}{\|v_1 - x_t\|^2} (v_1 - x_t)$  $r_2 = r_1 - \lambda_1 u_1$ 



• 
$$v_0 \in \arg \max_{v \in \mathcal{V}} \langle -\nabla f(x_t), v \rangle$$
  
 $\lambda_0 u_0 = \frac{\langle -\nabla f(x_t), v_0 - x_t \rangle}{\|v_0 - x_t\|^2} (v_0 - x_t)$   
 $r_1 = -\nabla f(x_t) - \lambda_0 u_0$ 

- $v_1 \in \arg \max_{v \in \mathcal{V}} \langle r_1, v \rangle$   $\lambda_1 u_1 = \frac{\langle r_1, v_1 - x_t \rangle}{\|v_1 - x_t\|^2} (v_1 - x_t)$  $r_2 = r_1 - \lambda_1 u_1$
- We could continue:
   v<sub>2</sub> ∈ arg max<sub>v∈V</sub> ⟨r<sub>2</sub>, v⟩



• 
$$v_0 \in \arg \max_{v \in \mathcal{V}} \langle -\nabla f(x_t), v \rangle$$
  
 $\lambda_0 u_0 = \frac{\langle -\nabla f(x_t), v_0 - x_t \rangle}{\|v_0 - x_t\|^2} (v_0 - x_t)$   
 $r_1 = -\nabla f(x_t) - \lambda_0 u_0$ 

- $v_1 \in \arg \max_{v \in \mathcal{V}} \langle r_1, v \rangle$   $\lambda_1 u_1 = \frac{\langle r_1, v_1 - x_t \rangle}{\|v_1 - x_t\|^2} (v_1 - x_t)$  $r_2 = r_1 - \lambda_1 u_1$
- We could continue:
   v<sub>2</sub> ∈ arg max<sub>v∈V</sub> ⟨r<sub>2</sub>, v⟩
- $d = \lambda_0 u_0 + \lambda_1 u_1$



• 
$$v_0 \in \arg \max_{v \in \mathcal{V}} \langle -\nabla f(x_t), v \rangle$$
  
 $\lambda_0 u_0 = \frac{\langle -\nabla f(x_t), v_0 - x_t \rangle}{\|v_0 - x_t\|^2} (v_0 - x_t)$   
 $r_1 = -\nabla f(x_t) - \lambda_0 u_0$ 

- $v_1 \in \arg \max_{v \in \mathcal{V}} \langle r_1, v \rangle$   $\lambda_1 u_1 = \frac{\langle r_1, v_1 - x_t \rangle}{\|v_1 - x_t\|^2} (v_1 - x_t)$  $r_2 = r_1 - \lambda_1 u_1$
- We could continue:
   v<sub>2</sub> ∈ arg max<sub>v∈V</sub> ⟨r<sub>2</sub>, v⟩
- $d = \lambda_0 u_0 + \lambda_1 u_1$
- $g_t = d/(\lambda_0 + \lambda_1)$



 How can we build a direction better aligned with −∇f(x<sub>t</sub>) and that allows to update x<sub>t+1</sub> without projection?

• 
$$v_0 \in \arg \max_{v \in \mathcal{V}} \langle -\nabla f(x_t), v \rangle$$
  
 $\lambda_0 u_0 = \frac{\langle -\nabla f(x_t), v_0 - x_t \rangle}{\|v_0 - x_t\|^2} (v_0 - x_t)$   
 $r_1 = -\nabla f(x_t) - \lambda_0 u_0$ 

- $v_1 \in \arg \max_{v \in \mathcal{V}} \langle r_1, v \rangle$   $\lambda_1 u_1 = \frac{\langle r_1, v_1 - x_t \rangle}{\|v_1 - x_t\|^2} (v_1 - x_t)$  $r_2 = r_1 - \lambda_1 u_1$
- We could continue:
   v<sub>2</sub> ∈ arg max<sub>v∈V</sub> ⟨r<sub>2</sub>, v⟩
- $d = \lambda_0 u_0 + \lambda_1 u_1$
- $g_t = d/(\lambda_0 + \lambda_1)$



The boosted direction g<sub>t</sub> is better aligned with −∇f(x<sub>t</sub>) than is the FW direction v<sub>0</sub> − x<sub>t</sub>

 How can we build a direction better aligned with −∇f(x<sub>t</sub>) and that allows to update x<sub>t+1</sub> without projection?

• 
$$v_0 \in \arg \max_{v \in \mathcal{V}} \langle -\nabla f(x_t), v \rangle$$
  
 $\lambda_0 u_0 = \frac{\langle -\nabla f(x_t), v_0 - x_t \rangle}{\|v_0 - x_t\|^2} (v_0 - x_t)$   
 $r_1 = -\nabla f(x_t) - \lambda_0 u_0$ 

- $v_1 \in \arg \max_{v \in \mathcal{V}} \langle r_1, v \rangle$   $\lambda_1 u_1 = \frac{\langle r_1, v_1 - x_t \rangle}{\|v_1 - x_t\|^2} (v_1 - x_t)$  $r_2 = r_1 - \lambda_1 u_1$
- We could continue:
   v<sub>2</sub> ∈ arg max<sub>v∈V</sub> ⟨r<sub>2</sub>, v⟩
- $d = \lambda_0 u_0 + \lambda_1 u_1$
- $g_t = d/(\lambda_0 + \lambda_1)$



The boosted direction g<sub>t</sub> is better aligned with -∇f(x<sub>t</sub>) than is the FW direction v<sub>0</sub> - x<sub>t</sub> and satisfies [x<sub>t</sub>, x<sub>t</sub> + g<sub>t</sub>] ⊆ C so we can update

$$x_{t+1} = x_t + \gamma_t g_t$$
 for any  $\gamma_t \in [0,1]$ 

Why  $[x_t, x_t + g_t] \subseteq C$ ? Let  $K_t$  be the number of alignment rounds. We have

$$d = \sum_{k=0}^{K_t-1} \lambda_k (v_k - x_t) \quad \text{where } \lambda_k > 0 \text{ and } v_k \in \mathcal{V}$$

Why  $[x_t, x_t + g_t] \subseteq \mathcal{C}$ ? Let  $K_t$  be the number of alignment rounds. We have

$$d = \sum_{k=0}^{K_t-1} \lambda_k (v_k - x_t)$$
 where  $\lambda_k > 0$  and  $v_k \in \mathcal{V}$ 

so if  $\Lambda_t = \sum_{k=0}^{K-1} \lambda_k$ , then

$$g_t = \frac{1}{\Lambda_t} \sum_{k=0}^{K_t-1} \lambda_k (v_k - x_t) = \underbrace{\left(\frac{1}{\Lambda_t} \sum_{k=0}^{K_t-1} \lambda_k v_k\right)}_{\in \mathcal{C}} - x_t$$

Why  $[x_t, x_t + g_t] \subseteq C$ ? Let  $K_t$  be the number of alignment rounds. We have

$$d = \sum_{k=0}^{K_t-1} \lambda_k (v_k - x_t)$$
 where  $\lambda_k > 0$  and  $v_k \in \mathcal{V}$ 

so if  $\Lambda_t = \sum_{k=0}^{K-1} \lambda_k$ , then

$$g_t = \frac{1}{\Lambda_t} \sum_{k=0}^{K_t-1} \lambda_k (v_k - x_t) = \underbrace{\left(\frac{1}{\Lambda_t} \sum_{k=0}^{K_t-1} \lambda_k v_k\right)}_{\in \mathcal{C}} - x_t$$

Thus,  $x_t + g_t \in \mathcal{C}$  so  $[x_t, x_t + g_t] \subseteq \mathcal{C}$  by convexity

**Algorithm** Finding a direction *g* well aligned with  $\nabla$  from a reference point *z* 

Inpu	ut: $z \in \mathcal{C}$ , $ abla \in \mathcal{H}$ , $K \in \mathbb{N} \setminus \{0\}$ , $\delta \in ]0, 1[$ .	
1:	$d_0 \leftarrow 0, \ \Lambda \leftarrow 0$	
2:	for $k = 0$ to $K - 1$ do	
3:	$r_k \leftarrow  abla - d_k$	⊳ <i>k</i> -th residual
4:	$v_k \leftarrow {\sf argmax}_{v \in \mathcal{V}} \langle r_k, v  angle$	▷ FW oracle
5:	$u_k \leftarrow rg\max_{u \in \{v_k - z, -d_k / \ d_k\ \}} \langle r_k, u \rangle$	
6:	$\lambda_k \leftarrow \langle r_k, u_k \rangle / \ u_k\ ^2$	
7:	$d'_k \leftarrow d_k + \lambda_k u_k$	
8:	if $align( abla, d'_k) - align( abla, d_k) \geqslant \delta$ then	
9:	$d_{k+1} \leftarrow d_k'$	
10:	$egin{aligned} & \Lambda_t \leftarrow egin{cases} \Lambda + \lambda_k &  ext{if } u_k = v_k - z \ & \Lambda(1 - \lambda_k / \ d_k\ ) &  ext{if } u_k = -d_k / \ d_k\  \end{aligned}$	
11:	else	
12:	break	⊳ exit <i>k</i> -loop
13:	$g \leftarrow d_k / \Lambda$	▷ normalization

**Algorithm** Finding a direction g well aligned with  $\nabla$  from a reference point z

Inpı	<b>it:</b> $z \in C$ , $\nabla \in H$ , $K \in \mathbb{N} \setminus \{0\}$ , $\delta \in ]0$ ,	L[.
1:	$d_0 \leftarrow 0, \ \Lambda \leftarrow 0$	
2:	for $k = 0$ to $K - 1$ do	
3:	$r_k \leftarrow  abla - d_k$	▷ k-th residual
4:	$m{v}_k \leftarrow {\sf argmax}_{m{v} \in \mathcal{V}} \langle m{r}_k,m{v}  angle$	▷ FW oracle
5:	$u_k \leftarrow \operatorname{argmax}_{u \in \{v_k - z, -d_k / \ d_k\ \}} \langle r_k, u_k \rangle$	ı>
6:	$\lambda_k \leftarrow \langle r_k, u_k \rangle / \ u_k\ ^2$	
7:	$d_k' \leftarrow d_k + \lambda_k u_k$	
8:	$if \; align(\nabla, d_k') - align(\nabla, d_k) \geqslant \delta \; t$	hen
9:	$d_{k+1} \leftarrow d_k'$	
10:	$\int \Lambda + \lambda_k$ if $u_k =$	$v_k - z$
	$\Lambda_t \leftarrow \left\{ \Lambda(1 - \lambda_k / \ d_k\ ) \text{ if } u_k = \right\}$	$= -d_k/\ d_k\ $
11:	else	
12:	break	⊳ exit <i>k</i> -loop
13:	$g \leftarrow d_k / \Lambda$	▷ normalization

• Technicality to ensure convergence of the procedure (Locatello et al., 2017)

**Algorithm** Finding a direction g well aligned with  $\nabla$  from a reference point z

Inpu	it: $z \in \mathcal{C}$ , $\nabla \in \mathcal{H}$ , $K \in \mathbb{N} \setminus \{0\}$ , $\delta \in ]0, 1[$ .	
1:	$d_0 \leftarrow 0,  \Lambda \leftarrow 0$	
2:	for $k = 0$ to $K - 1$ do	
3:	$r_k \leftarrow  abla - d_k$	⊳ <i>k</i> -th residual
4:	$v_k \leftarrow {\sf argmax}_{v \in \mathcal{V}} \langle r_k, v  angle$	▷ FW oracle
5:	$u_k \leftarrow rg\max_{u \in \{v_k - z, -d_k / \ d_k\ \}} \langle r_k, u  angle$	
6:	$\lambda_k \leftarrow \langle \mathbf{r}_k, \mathbf{u}_k \rangle / \ \mathbf{u}_k\ ^2$	
7:	$d_k' \leftarrow d_k + \lambda_k u_k$	
8:	if $align( abla, d'_k) - align( abla, d_k) \geqslant \delta$ then	
9:	$d_{k+1} \leftarrow d_k'$	
10.	$\int \Lambda + \lambda_k \qquad \text{if } u_k = v_k - z$	
10:	$\Lambda_t \leftarrow \left\{ \Lambda(1 - \lambda_k / \ d_k\ )  \text{if } u_k = -d_k / \ d_k\  \right\}$	
11:	else	
12:	break	⊳ exit <i>k</i> -loop
13:	$g \leftarrow d_k / \Lambda$	normalization

- Technicality to ensure convergence of the procedure (Locatello et al., 2017)
- The stopping criterion is an alignment improvement condition (typically  $\delta=10^{-3}$  and  $K=+\infty)$

AlgorithmFrank-Wolfe (FW)Input: $x_0 \in C, \ \gamma_t \in [0, 1].$ 1:for t = 0 to T - 1 do2: $v_t \leftarrow \arg\min(\nabla f(x_t), v)$ 3: $x_{t+1} \leftarrow x_t + \gamma_t(v_t - x_t)$ 

AlgorithmBoosted Frank-Wolfe (BoostFW)Input:  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ ,  $K \in \mathbb{N} \setminus \{0\}$ ,  $\delta \in ]0, 1[$ .1: for t = 0 to T - 1 do2:  $g_t \leftarrow \text{procedure}(x_t, -\nabla f(x_t), K, \delta)$ 

3: 
$$x_{t+1} \leftarrow x_t + \gamma_t g_t$$

AlgorithmFrank-Wolfe (FW)Input:  $x_0 \in C, \ \gamma_t \in [0, 1].$ 1: for t = 0 to T - 1 do2:  $v_t \leftarrow \underset{v \in \mathcal{V}}{\operatorname{arg\,min}} \langle \nabla f(x_t), v \rangle$ 3:  $x_{t+1} \leftarrow x_t + \gamma_t (v_t - x_t)$ 

**Algorithm** Boosted Frank-Wolfe (BoostFW)

Input:  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ ,  $K \in \mathbb{N} \setminus \{0\}$ ,  $\delta \in ]0, 1[$ .

- 1: for t = 0 to T 1 do
- 2:  $g_t \leftarrow \text{procedure}(x_t, -\nabla f(x_t), K, \delta)$
- 3:  $x_{t+1} \leftarrow x_t + \gamma_t g_t$

AlgorithmFrank-Wolfe (FW)Input:  $x_0 \in C, \ \gamma_t \in [0, 1].$ 1: for t = 0 to T - 1 do2:  $v_t \leftarrow \arg\min_{v \in \mathcal{V}} \langle \nabla f(x_t), v \rangle$ 3:  $x_{t+1} \leftarrow x_t + \gamma_t (v_t - x_t)$ 

Algorithm Boosted Frank-Wolfe (BoostFW)

Input:  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ ,  $K \in \mathbb{N} \setminus \{0\}$ ,  $\delta \in ]0, 1[$ .

1: for 
$$t = 0$$
 to  $T - 1$  do

2: 
$$g_t \leftarrow \text{procedure}(x_t, -\nabla f(x_t), K, \delta)$$

3: 
$$x_{t+1} \leftarrow x_t + \gamma_t g_t$$



AlgorithmBoosted Frank-Wolfe (BoostFW)Input:  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ ,  $K \in \mathbb{N} \setminus \{0\}$ ,  $\delta \in ]0, 1[$ .1: for t = 0 to T - 1 do2:  $g_t \leftarrow \text{procedure}(x_t, -\nabla f(x_t), K, \delta)$ 3:  $x_{t+1} \leftarrow x_t + \gamma_t g_t$ 







AlgorithmBoosted Frank-Wolfe (BoostFW)Input:  $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ ,  $K \in \mathbb{N} \setminus \{0\}$ ,  $\delta \in ]0, 1[$ .1: for t = 0 to T - 1 do2:  $g_t \leftarrow \text{procedure}(x_t, -\nabla f(x_t), K, \delta)$ 3:  $x_{t+1} \leftarrow x_t + \gamma_t g_t$ 



 $x^* = x_1$ 

What is the convergence rate of BoostFW?



AlgorithmBoosted Frank-Wolfe (BoostFW)Input: $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ ,  $K \in \mathbb{N} \setminus \{0\}$ ,  $\delta \in ]0, 1[$ .1:for t = 0 to T - 1 do2: $g_t \leftarrow \text{procedure}(x_t, -\nabla f(x_t), K, \delta)$ 3: $x_{t+1} \leftarrow x_t + \gamma_t g_t$ 



 $x^* = x_1$ 

- What is the convergence rate of BoostFW?
- Is BoostFW expensive in practice?



AlgorithmBoosted Frank-Wolfe (BoostFW)Input: $x_0 \in C$ ,  $\gamma_t \in [0, 1]$ ,  $K \in \mathbb{N} \setminus \{0\}$ ,  $\delta \in ]0, 1[$ .1:for t = 0 to T - 1 do2: $g_t \leftarrow \text{procedure}(x_t, -\nabla f(x_t), K, \delta)$ 3: $x_{t+1} \leftarrow x_t + \gamma_t g_t$ 





- What is the convergence rate of BoostFW?
- Is BoostFW expensive in practice?
- How does it compare to the state-of-the-art?

 Let N<sub>t</sub> be the number of iterations up to t where at least 2 rounds of alignment were performed (FW = always 1 round)

#### Theorem

Let  $C \subset \mathcal{H}$  be a compact convex set with diameter D and  $f : \mathcal{H} \to \mathbb{R}$  be a *L*-smooth, convex, and  $\mu$ -gradient dominated function, and let  $x_0 \in \arg\min_{v \in \mathcal{V}} \langle \nabla f(y), v \rangle$  for some  $y \in C$ . Set  $\gamma_t = \min\left\{\frac{\langle -\nabla f(x_t), g_t \rangle}{L \|g_t\|^2}, 1\right\}$  ("short step") and suppose that  $N_t \ge \omega t^p$  where  $p \in ]0, 1]$ . Then

$$f(x_t) - \min_{\mathcal{C}} f \leqslant \frac{LD^2}{2} \exp\left(-\delta^2 \frac{\mu}{L} \omega t^p\right)$$

 Let N<sub>t</sub> be the number of iterations up to t where at least 2 rounds of alignment were performed (FW = always 1 round)

#### Theorem

Let  $C \subset \mathcal{H}$  be a compact convex set with diameter D and  $f : \mathcal{H} \to \mathbb{R}$  be a *L*-smooth, convex, and  $\mu$ -gradient dominated function, and let  $x_0 \in \arg\min_{v \in \mathcal{V}} \langle \nabla f(y), v \rangle$  for some  $y \in C$ . Set  $\gamma_t = \min\left\{\frac{\langle -\nabla f(x_t), g_t \rangle}{L \|g_t\|^2}, 1\right\}$  ("short step") and suppose that  $N_t \ge \omega t^p$  where  $p \in [0, 1]$ . Then

$$f(x_t) - \min_{\mathcal{C}} f \leqslant \frac{LD^2}{2} \exp\left(-\delta^2 \frac{\mu}{L} \omega t^p\right)$$

• The assumption  $N_t \ge \omega t^p$  simply states that  $N_t$  is nonnegligeable, i.e., that the boosting procedure is active

 Let N<sub>t</sub> be the number of iterations up to t where at least 2 rounds of alignment were performed (FW = always 1 round)

#### Theorem

Let  $C \subset \mathcal{H}$  be a compact convex set with diameter D and  $f : \mathcal{H} \to \mathbb{R}$  be a *L*-smooth, convex, and  $\mu$ -gradient dominated function, and let  $x_0 \in \arg\min_{v \in \mathcal{V}} \langle \nabla f(y), v \rangle$  for some  $y \in C$ . Set  $\gamma_t = \min\left\{\frac{\langle -\nabla f(x_t), g_t \rangle}{L \|g_t\|^2}, 1\right\}$  ("short step") and suppose that  $N_t \ge \omega t^p$  where  $p \in [0, 1]$ . Then

$$f(x_t) - \min_{\mathcal{C}} f \leqslant \frac{LD^2}{2} \exp\left(-\delta^2 \frac{\mu}{L} \omega t^p\right)$$

- The assumption  $N_t \ge \omega t^{\rho}$  simply states that  $N_t$  is nonnegligeable, i.e., that the boosting procedure is active
- Else BoostFW reduces to FW and the convergence rate is  $\frac{4LD^2}{t+2}$
## Boosting Frank-Wolfe

 Let N<sub>t</sub> be the number of iterations up to t where at least 2 rounds of alignment were performed (FW = always 1 round)

#### Theorem

Let  $C \subset \mathcal{H}$  be a compact convex set with diameter D and  $f : \mathcal{H} \to \mathbb{R}$  be a *L*-smooth, convex, and  $\mu$ -gradient dominated function, and let  $x_0 \in \arg\min_{v \in \mathcal{V}} \langle \nabla f(y), v \rangle$  for some  $y \in C$ . Set  $\gamma_t = \min\left\{\frac{\langle -\nabla f(x_t), g_t \rangle}{L \|g_t\|^2}, 1\right\}$  ("short step") and suppose that  $N_t \ge \omega t^p$  where  $p \in [0, 1]$ . Then

$$f(x_t) - \min_{\mathcal{C}} f \leqslant \frac{LD^2}{2} \exp\left(-\delta^2 \frac{\mu}{L} \omega t^p\right)$$

- The assumption  $N_t \ge \omega t^{\rho}$  simply states that  $N_t$  is nonnegligeable, i.e., that the boosting procedure is active
- Else BoostFW reduces to FW and the convergence rate is  $\frac{4LD^2}{t+2}$
- In practice,  $N_t pprox t$  (so  $\omega \lesssim 1$  and p=1)

 We compare BoostFW to AFW, BCG, and DICG on a series of experiments involving various objective functions and feasible regions

 We compare BoostFW to AFW, BCG, and DICG on a series of experiments involving various objective functions and feasible regions

$$\begin{split} \min_{\substack{x \in \mathbb{R}^n \\ \textbf{s.t. } \|x\|_1 \leqslant \tau}} & \sum_{a \in \mathcal{A}} \tau_a x_a \left( 1 + 0.03 \left( \frac{x_a}{c_a} \right)^4 \right) \\ \text{s.t. } \|x\|_1 \leqslant \tau & \text{s.t. } \|x\|_1 \leqslant \tau \\ & \sum_{\substack{r \in \mathcal{R}_{i,j}}} y_r = d_{i,j} \quad (i,j) \in \mathcal{S} \\ & y_r \geqslant 0 \quad r \in \mathcal{R}_{i,j}, \ (i,j) \in \mathcal{S} \end{split}$$

$$\min_{x \in \mathbb{R}^n} \frac{1}{m} \sum_{i=1}^m \ln(1 + \exp(-y_i a_i^\top x))$$
s.t.  $||x||_1 \leq \tau$ 

$$\begin{split} \min_{X \in \mathbb{R}^{m \times n}} \frac{1}{|\mathcal{I}|} \sum_{(i,j) \in \mathcal{I}} h_{\rho}(Y_{i,j} - X_{i,j}) \\ \text{s.t. } \|X\|_{\text{nuc}} \leqslant \tau \end{split}$$

 We compare BoostFW to AFW, BCG, and DICG on a series of experiments involving various objective functions and feasible regions

$$\begin{split} \min_{x \in \mathbb{R}^{|\mathcal{A}|}} & \sum_{a \in \mathcal{A}} \tau_a x_a \left( 1 + 0.03 \left( \frac{x_a}{c_a} \right)^4 \right) \\ \min_{x \in \mathbb{R}^n} \|y - Ax\|_2^2 & \text{s.t. } x_a = \sum_{r \in \mathcal{R}} \mathbb{1}_{\{a \in r\}} y_r \quad a \in \mathcal{A} \\ & \sum_{r \in \mathcal{R}_{i,j}} y_r = d_{i,j} \quad (i,j) \in \mathcal{S} \\ & y_r \geqslant 0 \qquad r \in \mathcal{R}_{i,j}, (i,j) \in \mathcal{S} \end{split}$$

$$\min_{x \in \mathbb{R}^n} \frac{1}{m} \sum_{i=1}^m \ln(1 + \exp(-y_i a_i^\top x))$$

$$\min_{x \in \mathbb{R}^{m \times n}} \frac{1}{|\mathcal{I}|} \sum_{(i,j) \in \mathcal{I}} h_{\rho}(Y_{i,j} - X_{i,j})$$

$$\text{s.t. } \|x\|_1 \leq \tau$$

$$\text{s.t. } \|X\|_{\text{nuc}} \leq \tau$$

• For BoostFW and AFW we also run the line search-free variants (the "short step" strategy) and label them with an "L"





### • Traffic assignment



- Sparse logistic regression on the Gisette dataset
- Collaborative filtering on the MovieLens 100k dataset



40

- DICG is known to perform particularly well on the video co-localization experiment (YouTube-Objects dataset)
- BoostDICG: application of our method to DICG



- DICG is known to perform particularly well on the video co-localization experiment (YouTube-Objects dataset)
- BoostDICG: application of our method to DICG



• (details)

DICG

$$egin{aligned} & a_t \leftarrow ext{away vertex} \ & v_t \leftarrow rgmin_{v \in \mathcal{V}} \langle 
abla f(x_t), v 
angle \ & x_{t+1} \leftarrow x_t + \gamma_t (v_t - a_t) \end{aligned}$$

#### BoostDICG

 $a_t \leftarrow \text{away vertex}$   $g_t \leftarrow \text{procedure}(a_t, -\nabla f(x_t), K, \delta)$  $x_{t+1} \leftarrow x_t + \gamma_t g_t$ 

- DICG is known to perform particularly well on the video co-localization experiment (YouTube-Objects dataset)
- BoostDICG: application of our method to DICG



• (details)

DICG

$$a_t \leftarrow \text{away vertex}$$

$$v_t \leftarrow \underset{v \in \mathcal{V}}{\arg\min} \langle \nabla f(x_t), v \rangle$$

$$x_{t+1} \leftarrow x_t + \gamma_t (v_t - a_t)$$

#### BoostDICG

 $a_t \leftarrow \text{away vertex}$   $g_t \leftarrow \text{procedure}(a_t, -\nabla f(x_t), K, \delta)$  $x_{t+1} \leftarrow x_t + \gamma_t g_t$ 

- DICG is known to perform particularly well on the video co-localization experiment (YouTube-Objects dataset)
- BoostDICG: application of our method to DICG



• (details)

DICG

$$a_t \leftarrow \text{away vertex}$$

$$v_t \leftarrow \underset{v \in \mathcal{V}}{\arg\min} \langle \nabla f(x_t), v \rangle$$

$$x_{t+1} \leftarrow x_t + \gamma_t (v_t - a_t)$$

#### BoostDICG

 $a_t \leftarrow \text{away vertex}$   $g_t \leftarrow \text{procedure}(a_t, -\nabla f(x_t), K, \delta)$  $x_{t+1} \leftarrow x_t + \gamma_t g_t$ 

• Projection-free algorithms are of considerable interest in optimization

- Projection-free algorithms are of considerable interest in optimization
- We have proposed an intuitive and generic boosting procedure to speed up Frank-Wolfe algorithms

- Projection-free algorithms are of considerable interest in optimization
- We have proposed an intuitive and generic boosting procedure to speed up Frank-Wolfe algorithms
- Although our method may perform more linear minimizations per iteration, the progress obtained greatly overcomes their cost

- Projection-free algorithms are of considerable interest in optimization
- We have proposed an intuitive and generic boosting procedure to speed up Frank-Wolfe algorithms
- Although our method may perform more linear minimizations per iteration, the progress obtained greatly overcomes their cost
- We focused on smooth convex objective functions, but we expect our method to provide significant gains in performance in other areas of optimization as well

- Projection-free algorithms are of considerable interest in optimization
- We have proposed an intuitive and generic boosting procedure to speed up Frank-Wolfe algorithms
- Although our method may perform more linear minimizations per iteration, the progress obtained greatly overcomes their cost
- We focused on smooth convex objective functions, but we expect our method to provide significant gains in performance in other areas of optimization as well
  - E.g., large-scale finite-sum/stochastic constrained optimization:

$$g_t \leftarrow \text{procedure}(x_t, -\tilde{\nabla}f(x_t), K, \delta)$$
$$x_{t+1} \leftarrow x_t + \gamma_t g_t$$

### References

- G. Braun, S. Pokutta, D. Tu, and S. Wright. Blended conditional gradients: the unconditioning of conditional gradients. *ICML*, 2019.
- M. D. Canon and C. D. Cullum. A tight upper bound on the rate of convergence of Frank-Wolfe algorithm. SIAM J. Control, 1968.

#### C. W. Combettes and S. Pokutta. Boosting Frank-Wolfe by chasing gradients. ICML, 2020.

- M. Frank and P. Wolfe. An algorithm for quadratic programming. Naval Res. Logist. Q., 1956.
- D. Garber and O. Meshi. Linear-memory and decomposition-invariant linearly convergent conditional gradient algorithm for structured polytopes. *NIPS*, 2016.
- M. Jaggi. Revisiting Frank-Wolfe: Projection-free sparse convex optimization. ICML, 2013.
- S. Lacoste-Julien and M. Jaggi. On the global linear convergence of Frank-Wolfe optimization variants. *NIPS*, 2015.
- G. Lan. The complexity of large-scale convex programming under a linear optimization oracle. Technical report, University of Florida, 2013.
- E. S. Levitin and B. T. Polyak. Constrained minimization methods. USSR Comp. Math. Math. Phys., 1966.
- F. Locatello, M. Tschannen, G. Rätsch, and M. Jaggi. Greedy algorithms for cone constrained optimization with convergence guarantees. *NIPS*, 2017.
- P. Wolfe. Convergence theory in nonlinear programming. *Integer and Nonlinear Programming*. North-Holland, 1970.