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Projection-free Distributed Online Convex Optimization with $O(\sqrt{T})$ Communication Complexity

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http://www.lambda.nju.edu.cn/wanyy Projection-free Distributed Online Learning

Introduction

- Background
- The Problem and Our Contributions

Our Algorithms

- D-BOCG for Full Information Setting
- D-BBCG for Bandit Setting

3 Experiments





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- Formal definition
- 1: for t = 1, 2, ..., T do
- 2: for each local learner $i \in [n]$ do
- 3: pick a decision $\mathbf{x}_i(t) \in \mathcal{K}$ receive a convex loss function $f_{t,i}(\mathbf{x}) : \mathcal{K} \to \mathbb{R}$
- 4: communicate with its neighbors and update $\mathbf{x}_i(t)$
- 5: end for
- 6: end for



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- the network is defined as G = (V, E), V = [n]
- each node $i \in [n]$ is a local learner
- node i can only communicate with its immediate neighbors

$$N_i = \{j \in V | (i,j) \in E\}$$

• the global loss function is defined as $f_t(\mathbf{x}) = \sum_{j=1}^n f_{t,j}(\mathbf{x})$



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- Regret of local learner i

$$R_{T,i} = \sum_{t=1}^{T} f_t(\mathbf{x}_i(t)) - \min_{\mathbf{x} \in \mathcal{K}} \sum_{t=1}^{T} f_t(\mathbf{x})$$



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Applications

- multi-agent coordination
- distributed tracking in sensor networks



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Projection-based Methods

Distributed Online Dual Averaging [Hosseini et al., 2013]

- 1: for each local learner $i \in [n]$ do
- 2: Play $\mathbf{x}_i(t)$ and compute $\mathbf{g}_i(t) = \nabla f_{t,i}(\mathbf{x}_i(t))$
- 3: $\mathbf{z}_i(t+1) = \sum_{j \in N_i} P_{ij} \mathbf{z}_j(t) + \mathbf{g}_i(t)$
- 4: $\mathbf{x}_i(t+1) = \Pi^{\psi}_{\mathcal{K}}(\mathbf{z}_i(t+1), \alpha(t))$

5: end for

- *P_{ij}* > 0 only if (*i*, *j*) ∈ *E* or *P_{ij}* = 0
- $\psi(\mathbf{x}) : \mathcal{K} \to \mathbb{R}$ is a proximal function, e.g., $\psi(\mathbf{x}) = \|\mathbf{x}\|_2^2$
- projection step: $\Pi^{\psi}_{\mathcal{K}}(\mathbf{z}, \alpha) = \operatorname{argmin}_{\mathbf{x} \in \mathcal{K}} \mathbf{z}^{\top} \mathbf{x} + \frac{1}{\alpha} \psi(\mathbf{x})$

•
$$\alpha(t) = O(1/\sqrt{t}) \rightarrow R_{T,i} = O(\sqrt{T})$$



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•
$$\alpha(t) = O(1/\sqrt{t}) \rightarrow R_{T,i} = O(\sqrt{T})$$

- Distributed Online Gradient Descent [Ram et al., 2010]
 - also need a projection step



Projection-free Methods

Motivation: the projection step could be time-consuming

 $\bullet\,$ if ${\cal K}$ is a trace norm ball, it requires SVD of a matrix



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Projection-free Methods

- Motivation: the projection step could be time-consuming
 - if \mathcal{K} is a trace norm ball, it requires SVD of a matrix
- Distributed Online Conditional Gradient [Zhang et al., 2017]
- 1: for each local learner $i \in [n]$ do
- 2: Play $\mathbf{x}_i(t)$ and compute $\mathbf{g}_i(t) = \nabla f_{t,i}(\mathbf{x}_i(t))$
- 3: $\mathbf{v}_i = \operatorname{argmin}_{\mathbf{x} \in \mathcal{K}} \nabla F_{t,i}(\mathbf{x}_i(t))^\top \mathbf{x}$
- 4: $\mathbf{x}_i(t+1) = \mathbf{x}_i(t) + \mathbf{s}_t(\mathbf{v}_i \mathbf{x}_i(t))$
- 5: $\mathbf{z}_i(t+1) = \sum_{j \in N_i} P_{ij}\mathbf{z}_j(t) + \mathbf{g}_i(t)$

6: **end for**

•
$$F_{t,i}(\mathbf{x}) = \eta \mathbf{z}_i(t)^\top \mathbf{x} + \|\mathbf{x} - \mathbf{x}_1(1)\|_2^2$$

•
$$\eta = O(T^{-3/4}), s_t = 1/\sqrt{t} \to R_{T,i} = O(T^{3/4})$$

- only contain linear optimization step (step 3)
- T communication rounds





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Question

Can the O(T) communication complexity of distributed online conditional gradient (D-OCG) be reduced?

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An affirmative and non-trivial answer

- distributed block online conditional gradient (D-BOCG)
- communication complexity: from O(T) to $O(\sqrt{T})$
- regret bound: $O(T^{3/4})$

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Question

Can the O(T) communication complexity of distributed online conditional gradient (D-OCG) be reduced?

An affirmative and non-trivial answer

- distributed block online conditional gradient (D-BOCG)
- communication complexity: from O(T) to $O(\sqrt{T})$
- regret bound: $O(T^{3/4})$
- An extension to the bandit setting
 - distributed block bandit conditional gradient (D-BBCG)
 - communication complexity: $O(\sqrt{T})$
 - high-probability regret bound: $\widetilde{O}(T^{3/4})$

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Our Algorithms

D-BOCG for Full Information Setting



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Main Idea

Delayed update mechanism

- block 1 ... block m ... block \sqrt{T}
- only update in the beginning of each block
- only need \sqrt{T} communication rounds



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Delayed update mechanism

- block 1 ... block m ... block \sqrt{T}
- only update in the beginning of each block
- only need \sqrt{T} communication rounds
- Iterative linear optimization steps
 - recall the update rules of D-OCG

$$\mathbf{v}_{i} = \underset{\mathbf{x} \in \mathcal{K}}{\operatorname{argmin}} \nabla F_{t,i}(\mathbf{x}_{i}(t))^{\top} \mathbf{x}$$
$$\mathbf{x}_{i}(t+1) = \mathbf{x}_{i}(t) + s_{t}(\mathbf{v}_{i} - \mathbf{x}_{i}(t))$$

- delayed update + D-OCG: a worse regret bound
- multiple linear optimization steps for each update



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Conditional Gradient with Stopping Condition (CGSC)

- CGSC [Garber and Kretzu, 2019]
- 1: Input: feasible set \mathcal{K} , $\epsilon > 0$, L, $F(\mathbf{x})$, \mathbf{x}_{in}
- 2: $\tau = 0, c_1 = x_{in}$
- 3: repeat
- 4: $\tau = \tau + 1$
- 5: $\mathbf{V}_{\tau} \in \underset{\mathbf{x} \in \mathcal{K}}{\operatorname{argmin}} \nabla F(\mathbf{C}_{\tau})^{\top} \mathbf{X}$ 6: $\mathbf{s} = \underset{\tau}{\operatorname{argmin}} F(\mathbf{C}_{\tau} + \mathbf{s})^{\top} \mathbf{X}$
- 6: $s_{\tau} = \underset{s \in [0,1]}{\operatorname{argmin}} F(\mathbf{c}_{\tau} + s(\mathbf{v}_{\tau} \mathbf{c}_{\tau}))$
- 7: $\mathbf{c}_{\tau+1} = \mathbf{c}_{\tau} + s_{\tau}(\mathbf{v}_{\tau} \mathbf{c}_{\tau})$
- 8: until $\nabla F(\mathbf{c}_{\tau})^{\top}(\mathbf{c}_{\tau}-\mathbf{v}_{\tau}) \leq \epsilon$ or $\tau = L$
- 9: return $\mathbf{x}_{out} = \mathbf{c}_{\tau}$
 - $F(\mathbf{x}_{out})$ is very small with appropriate L and ϵ
 - it was widely studied [Frank and Wolfe, 1956, Jaggi, 2013]



The Proposed D-BOCG Algorithm

1: Initialization: choose $\{\mathbf{x}_i(1) = \mathbf{0} \in \mathcal{K} | i \in V\}$ and set $\{z_i(1) = 0 | i \in V\}$ 2: for $t = 1, \dots, T$ do $m_t = \lfloor t/K \rfloor$ 3: for each local learner $i \in V$ do 4: if t > 1 and mod(t, K) = 1 then 5: $\widehat{\mathbf{g}}_i(m_t-1) = \sum_{k=t-K}^{t-1} \mathbf{g}_i(k)$ 6: $\mathbf{z}_i(m_t) = \sum_{i \in N_i} P_{ii} \mathbf{z}_i(m_t - 1) + \widehat{\mathbf{g}}_i(m_t - 1)$ 7: define $F_{m_t i}(\mathbf{x}) = \eta \mathbf{z}_i (m_t)^\top \mathbf{x} + \|\mathbf{x}\|_2^2$ 8: $\mathbf{x}_i(m_t) = \text{CGSC}(\mathcal{K}, \epsilon, L, F_{m_t}(\mathbf{x}), \mathbf{x}_i(m_t - 1))$ 9: end if 10: play $\mathbf{x}_i(m_t)$ and observe $\mathbf{g}_i(t) = \nabla f_{t,i}(\mathbf{x}_i(m_t))$ 11: end for 12: 13: end for



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Regret of D-BOCG

Theorem 1

Let $\eta = O(T^{-3/4})$, $\epsilon = O(T^{-1/2})$, $K = \sqrt{T}$ and $L = O(\sqrt{T})$. For any $i \in V$, D-BOCG has

$$R_{T,i} \leq O(GRT^{3/4}).$$

Assumptions

•
$$|f_{t,i}(\mathbf{x}) - f_{t,i}(\mathbf{y})| \le G \|\mathbf{x} - \mathbf{y}\|_2$$
 for any $\mathbf{x}, \mathbf{y} \in \mathcal{K}$

- $r\mathcal{B}^d \subseteq \mathcal{K} \subseteq R\mathcal{B}^d$, \mathcal{B}^d is the unit Euclidean ball
- $P \in \mathbb{R}^{n \times n}$ is symmetric and doubly stochastic, i.e.,

$$P = P^{\top}, \mathbf{1}^{\top}P = \mathbf{1}^{\top}, P\mathbf{1} = \mathbf{1}$$



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Remarks

- regret bound: $R_{T,i} = O(T^{3/4})$
- #communication rounds: $T/K = \sqrt{T}$
- #linear optimization steps: LT/K = O(T)



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Standard Technique

Bandit setting

- only the loss value is available to learners
- the main challenge is due to the lack of gradient



Standard Technique

- Bandit setting
 - only the loss value is available to learners
 - the main challenge is due to the lack of gradient
- One-point Gradient Estimator [Flaxman et al., 2005]
 - δ-smoothed version of f(x)

$$\widehat{f}_{\delta}(\mathbf{x}) = \mathbb{E}_{\mathbf{u} \sim \mathcal{B}^d}[f(\mathbf{x} + \delta \mathbf{u})]$$

• let $\delta > 0$ and S^d be the unit sphere

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• only observe the single value $f(\mathbf{x} + \delta \mathbf{u})$



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ight]$$

- only observe the single value $f(\mathbf{x} + \delta \mathbf{u})$
- A smaller set $\mathcal{K}_{\delta} \subseteq \mathcal{K}$

•
$$\mathcal{K}_{\delta} = (1 - \delta/r)\mathcal{K} = \{(1 - \delta/r)\mathbf{x} | \mathbf{x} \in \mathcal{K}\}, \, \mathbf{0} < \delta \leq r$$

• $\mathbf{x} + \delta \mathbf{u} \in \mathcal{K}$ for $\mathbf{x} \in \mathcal{K}_{\delta}, \mathbf{u} \sim \mathcal{S}$



The Proposed D-BBCG Algorithm

1: Initialization: choose $\{\mathbf{x}_i(1) = \mathbf{0} \in \mathcal{K}_{\delta} | i \in V\}$ and set $\{z_i(1) = 0 | i \in V\}$ 2: for $t = 1, \dots, T$ do $m_t = \lceil t/K \rceil$ 3: for each local learner $i \in V$ do 4. 5: if t > 1 and mod(t, K) = 1 then $\widehat{\mathbf{q}}_i(m_t-1) = \sum_{k=t}^{t-1} \kappa \mathbf{q}_i(k)$ 6: $\mathbf{z}_i(m_t) = \sum_{i \in N_i} P_{ii} \mathbf{z}_i(m_t - 1) + \widehat{\mathbf{g}}_i(m_t - 1)$ 7: define $F_{m_t,i}(\mathbf{x}) = \eta \mathbf{z}_i(m_t)^\top \mathbf{x} + \|\mathbf{x}\|_2^2$ 8: $\mathbf{x}_i(m_t) = \text{CGSC}(\mathcal{K}_{\delta}, \epsilon, L, F_{m_t,i}(\mathbf{x}), \mathbf{x}_i(m_t - 1))$ 9: 10: end if $\mathbf{u}_i(t) \sim \mathcal{S}^d$ 11: play $\mathbf{y}_i(t) = \mathbf{x}_i(m_t) + \delta \mathbf{u}_i(t)$ and observe $f_{t,i}(\mathbf{y}_i(t))$ 12: $\mathbf{g}_i(t) = \frac{d}{s} f_{t,i}(\mathbf{y}_i(t)) \mathbf{u}_i(t)$ 13: end for 14: 15: end for ・ 回 ト ・ ヨ ト ・ ヨ ト





Regret of D-BBCG

Theorem 2

Let
$$\eta = O(T^{-3/4})$$
, $\delta = O(T^{-1/4})$, $\epsilon = O(T^{-1/2})$, $K = T^{1/2}$ and $L = O(\sqrt{T})$. For any $i \in V$, with high probability, D-BBCG has $R_{T,i} \leq \widetilde{O}(T^{3/4})$.

Additional Assumption

all local loss functions are chosen beforehand



Regret of D-BBCG

Theorem 2

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Additional Assumption

- all local loss functions are chosen beforehand
- Remarks
 - high-probability regret bound: $R_{T,i} = \widetilde{O}(T^{3/4})$
 - #communication rounds: $T/K = \sqrt{T}$
 - #linear optimization steps: LT/K = O(T)



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Experimental Settings

Distributed multiclass classification [Zhang et al., 2017]

- 1: for t = 1, 2, ..., T do
- 2: for each local learner $i \in [n]$ do
- 3: receive an example $\mathbf{e}_i(t) \in \mathbb{R}^k$, and choose $X_i(t) = [\mathbf{x}_1^\top; \mathbf{x}_2^\top; \cdots; \mathbf{x}_h^\top] \in \mathcal{K}$
- 4: receive the true label $y_i(t)$, and suffer the multivariate logistic loss

$$f_{t,i}(X_i(t)) = \log\left(1 + \sum_{\ell \neq y_i(t)} e^{\mathbf{x}_\ell^\top \mathbf{e}_i(t) - \mathbf{x}_{y_i(t)}^\top \mathbf{e}_i(t)}\right)$$

- 5: communicate with its neighbors and update $X_i(t)$
- 6: end for
- 7: end for
 - $\mathcal{K} = \{X \in \mathbb{R}^{h \times k} | \|X\|_* \le \tau\}$, where $\|X\|_*$ denotes the trace norm of X and τ is a constant
 - the network is a cycle graph with 9 nodes



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Experimental Results

aloi dataset from the LIBSVM repository [Chang and Lin, 2011]



Experimental Results

poker dataset from the LIBSVM repository



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Conclusion and Future Work

Conclusion

- D-BOCG enjoys an $O(T^{3/4})$ regret bound with only $O(\sqrt{T})$ communication rounds
- D-BBCG for bandit setting enjoys a high-probability $\widetilde{O}(T^{3/4})$ regret bound with only $O(\sqrt{T})$ communication rounds

Future Work

 improve the regret bound of projection-free distributed online learning by utilizing the curvature of functions



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Thanks!



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